

# A set-theoretical investigation of Pāṇini's Śivasūtras

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# Pāṇini's Śivasūtras

*anubandha*

*sūtras*

1.	a	i	u			N
2.				r	l	K
3.		e	o			Ñ
4.		ai	au			C
5.	h	y	v	r		T
6.					l	N
7.	ñ	m	ṅ	ṇ	n	M
8.	jh	bh				Ñ
9.			gh	ḍh	dh	S
10.	j	b	g	ḍ	d	S
11.	kh	ph	ch	ṭh	th	
			c	ṭ	t	V
12.	k	p				Y
13.		ś	ṣ	s		R
14.	h					L

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*anubandha*

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1.	a	i	u			N̄
2.				r̄	l̄	K̄
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5.	h	y	v	r		T̄
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			c	ṭ	t	V
12.	k	p				Y
13.		ś	ṣ	s		R
14.	h					L

# Phonological classes/ *pratyāhāras*

1.	a	i	u				Ṇ
2.				ṛ	ḷ		Ḳ
3.		e	o				Ṇ̇
4.		ai	au				C
5.	h	y	v	r			Ṛ
6.						l	Ṇ
7.	ñ	m	ṅ	ṇ	n		M
8.	jh	bh					Ṇ̃

Phonological classes are denoted by *pratyāhāras*.

E.g., the *pratyāhāra* *iC* denotes the set of segments in the continuous sequence starting with *i* and ending with *au*, the last element before the *anubandha* *C*.

# Phonological classes/ *pratyāhāras*

1.	a	⓪	u				Ṇ
2.				ṛ	ḷ		Ḳ
3.		e	o				Ṇ̇
4.		ai	au				Ḳ̇
5.	h	y	v	r			Ṫ
6.						l	Ṇ
7.	ñ	m	ṅ	ṇ	n		Ṣ
8.	jh	bh					Ṇ̇

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1.	a	<b>i</b>	u					N
2.				<b>ṛ</b>	<b>ḷ</b>			K
3.		<b>e</b>	<b>o</b>					Ñ
4.		<b>ai</b>	<b>au</b>					<b>C</b>
5.	h	y	v	r				Ṭ
6.						l		N
7.	ñ	m	ṅ	ṇ	n			M
8.	jh	bh						Ñ̃

*iC*

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# Questions around Pāṇini's Śivasūtras

- Are Pāṇini's Śivasūtras perfect?
  - Is the number of stop markers minimal?
  - Is the duplication of 'h' necessary?
- Is it possible to decide whether a set of sets has a Śivasūtras-style representation?
- How to construct an optimal Śivasūtras-style representation?

# Basic concepts

S-encodable set of sets:  $\Phi = \{\{d,e\}, \{b,c,d,f,g,h,i\}, \{a,b\}, \{f,i\}, \{c,d,e,f,g,h,i\}, \{g,h\}\}$



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S-alphabet  $(\mathcal{A}, \Sigma, <)$  of  $\Phi$ :

e d  $M_1$  c i f  $M_2$  g h  $M_3$  b  $M_4$  a  $M_5$

alphabet

marker

total order on  $\mathcal{A} \cup \Sigma$

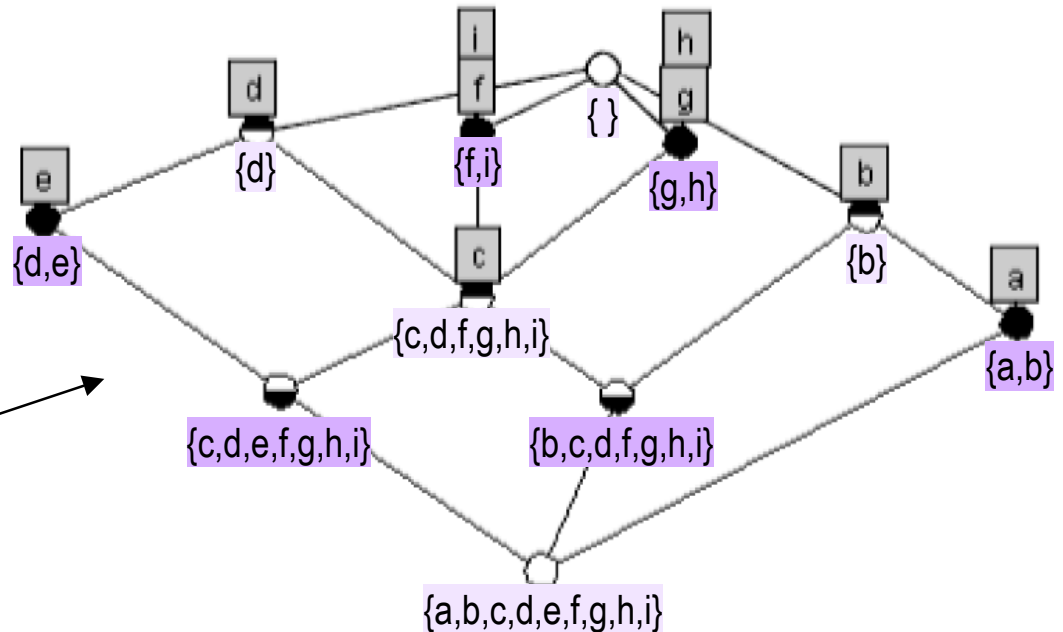
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alphabet  $\nearrow$   
 marker  $\nearrow$   
 total order on  $\mathcal{A} \cup \Sigma$   $\nearrow$



Hasse-diagram  $(\mathcal{H}(\Phi), \supseteq)$   
 for the set of intersections  
 of elements from  $\Phi \cup \{\mathcal{A}\}$ :

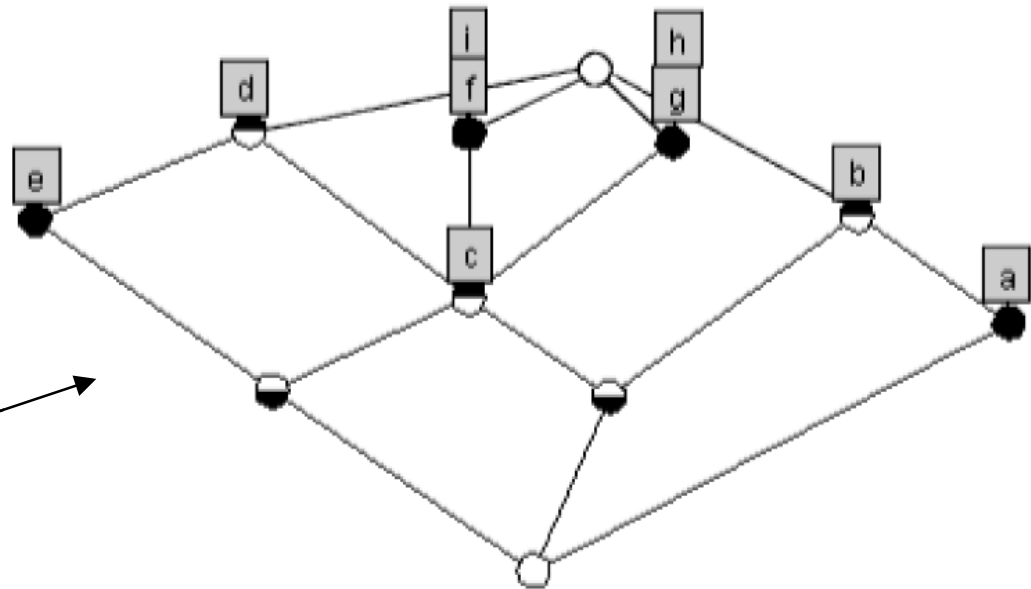
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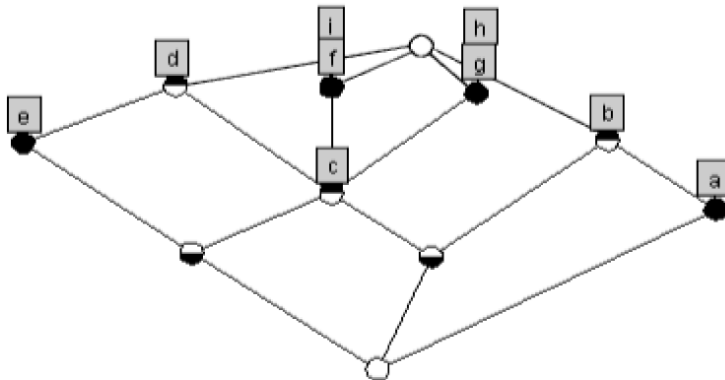
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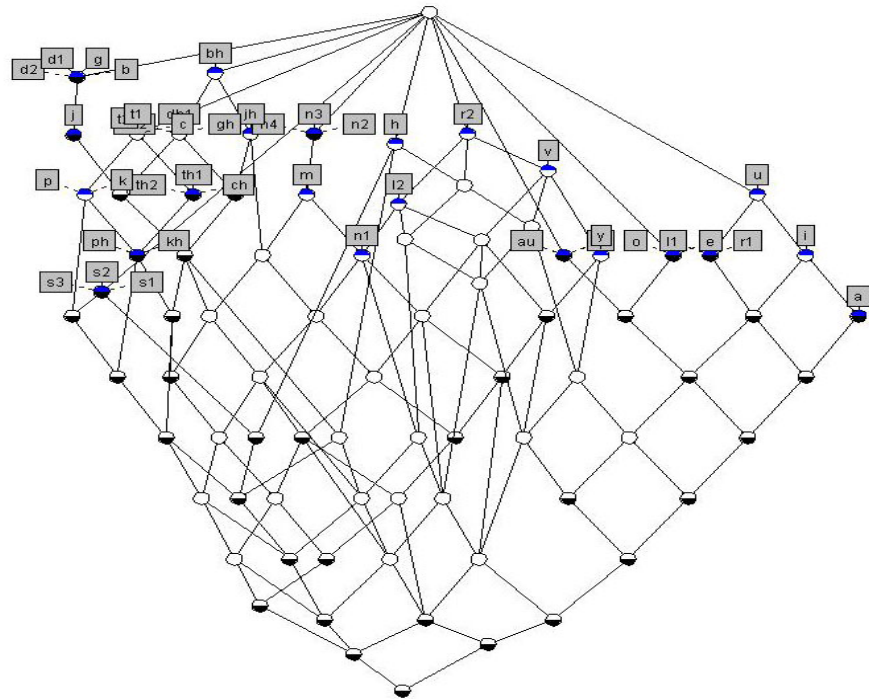
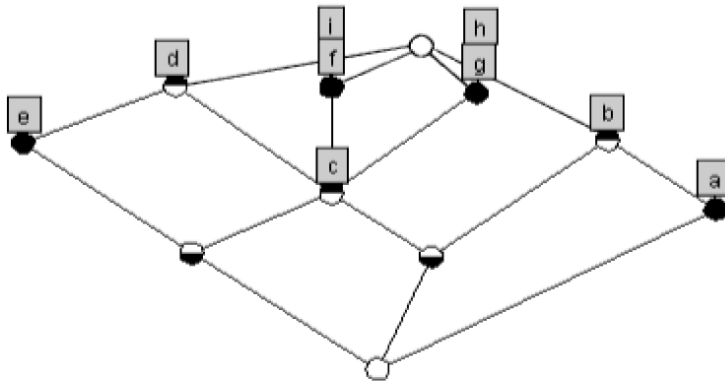
# S-encodability and planar Hasse-diagrams

If  $\Phi$  is S-encodable, then the Hasse-diagram of  $(\mathcal{H}(\Phi), \supseteq)$  is planar



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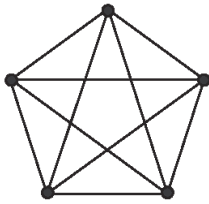


Hasse-diagram for Pāṇini's *pratyāhāras*

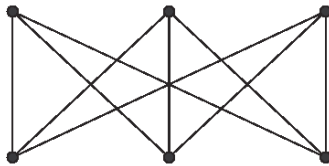
# S-encodability and planar Hasse-diagrams

*Criterion of Kuratowski:*

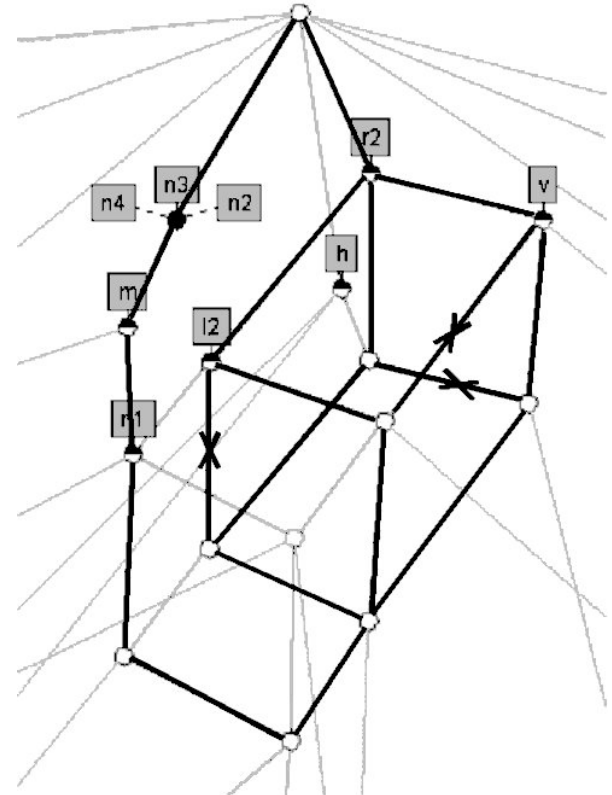
A graph is planar iff it has neither  $K_5$  nor  $K_{3,3}$  as a *minor*.



$K_5$



$K_{3,3}$

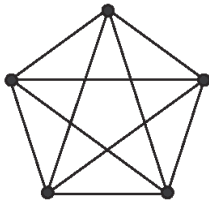


part of the Hasse-diagram for Pāṇini's pratyāhāras

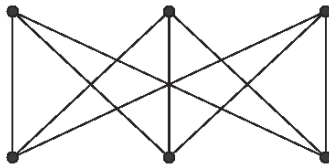
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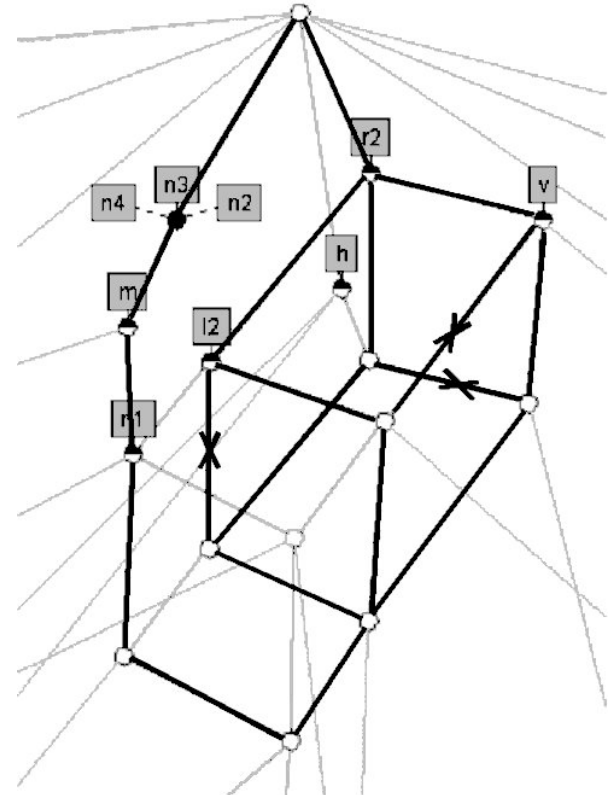
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$K_5$

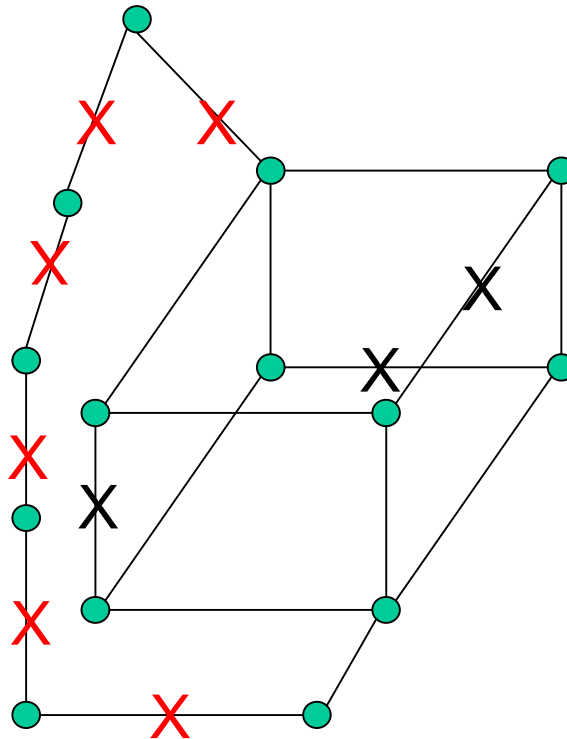


$K_{3,3}$



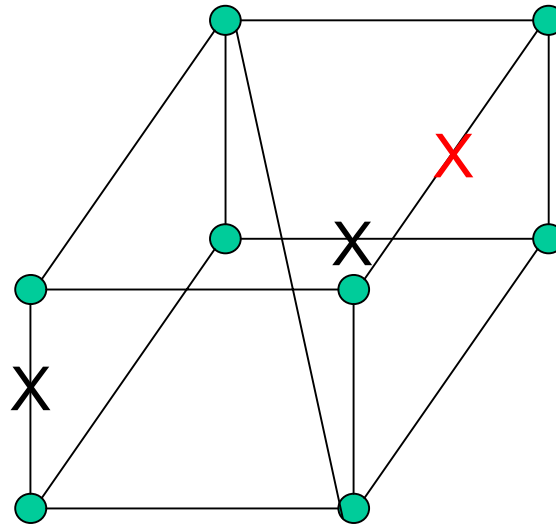
part of the Hasse-diagram for Pāṇini's pratyāhāras

# K5 is a minor of the Hasse-diagram for the pratyāhāras

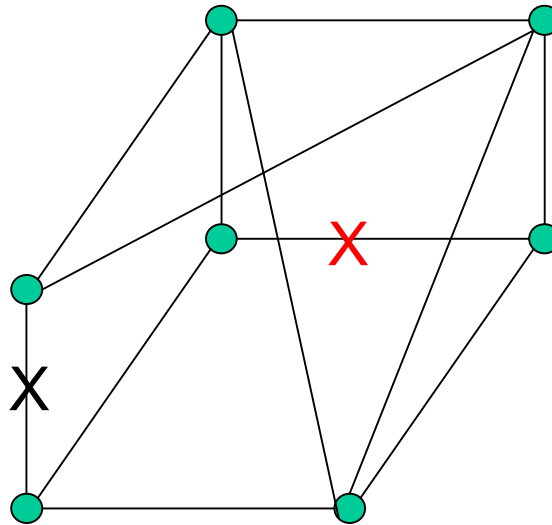




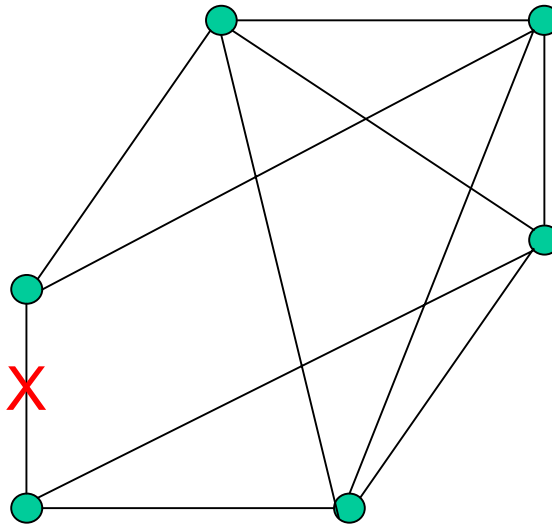
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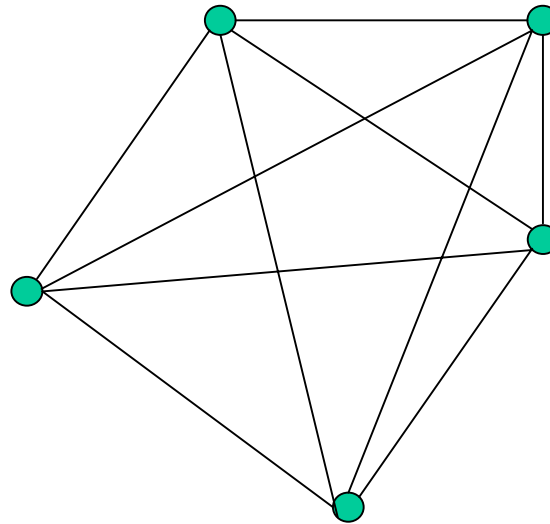
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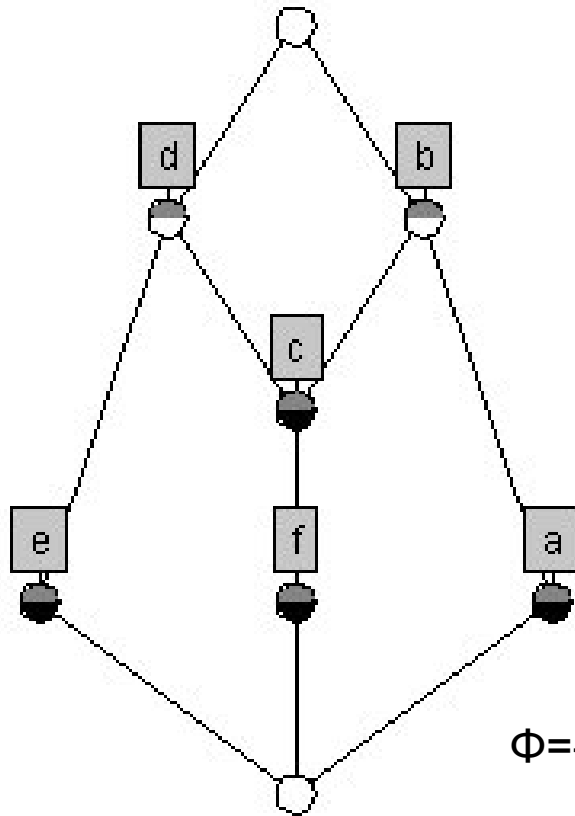
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# K5 is a minor of the Hasse-diagram for the pratyāhāras



# We are not done yet!



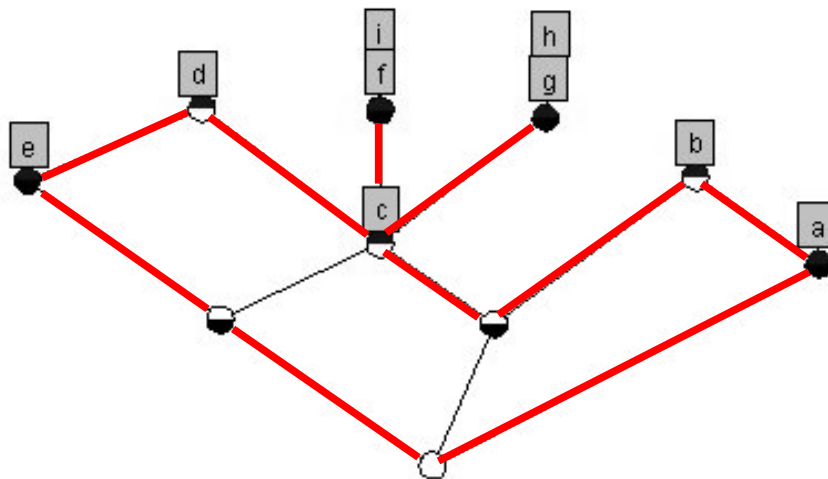
plane but not S-encodable!

$$\Phi = \{\{d,e\}, \{b,c,d,f\}, \{a,b\}, \{b,c,d\}\}$$

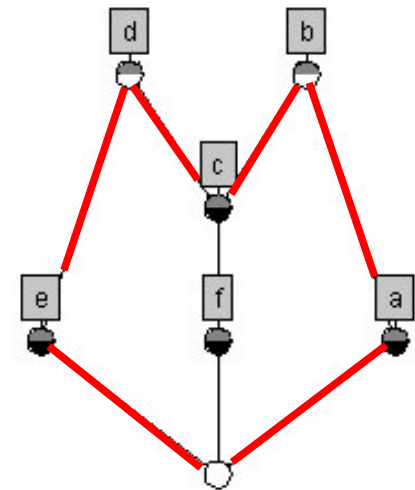
# Existence of S-alphabets

An S-alphabet of  $\Phi$  exists iff

1. a plane Hasse-diagram for  $(\mathcal{H}(\Phi), \supseteq)$  exists and
2. for every  $b \in \mathcal{A}$  the S-graph contains a node labeled with  $b$

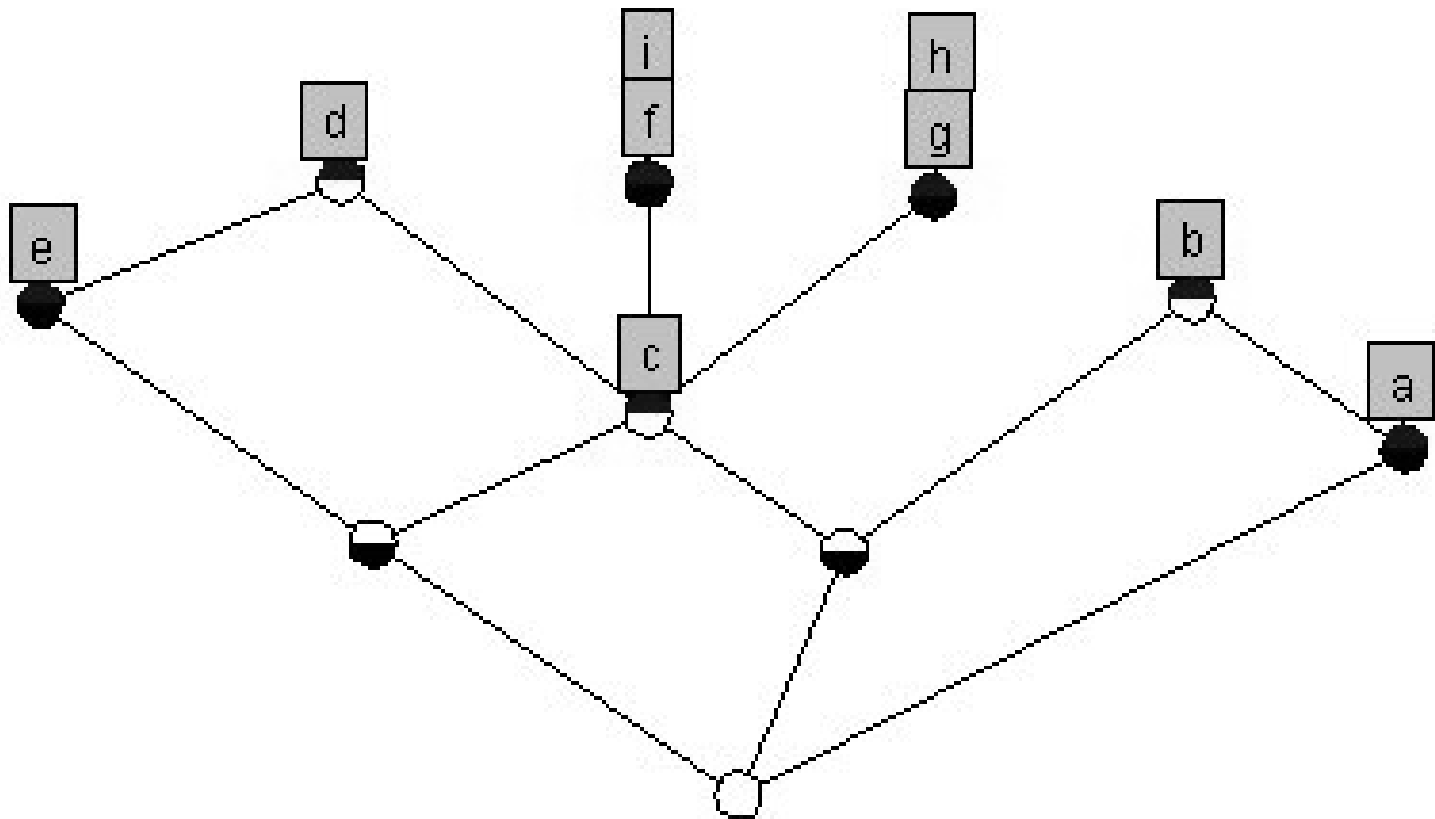


S-encodable



not S-encodable

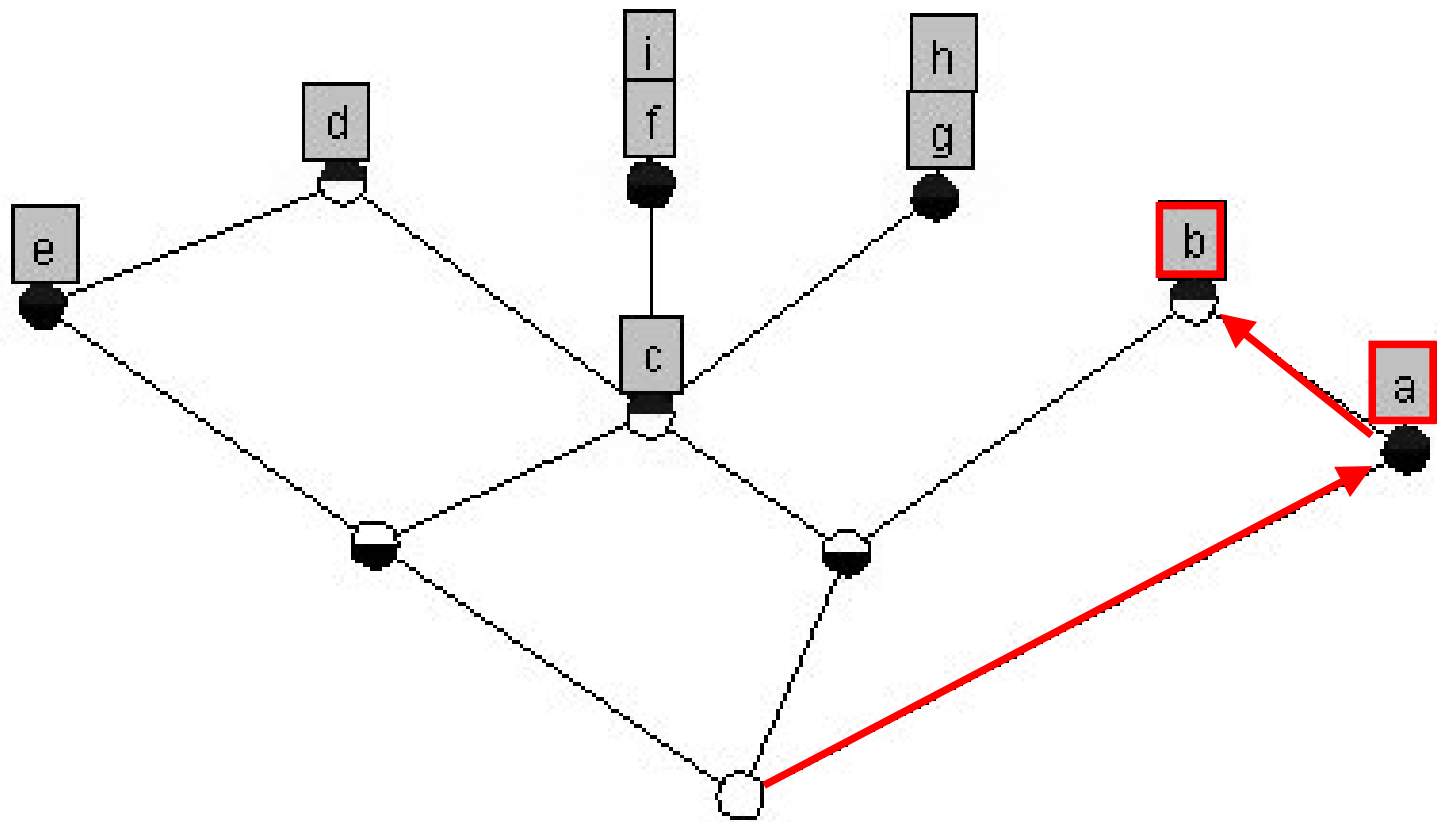
# Construction of S-alphabets





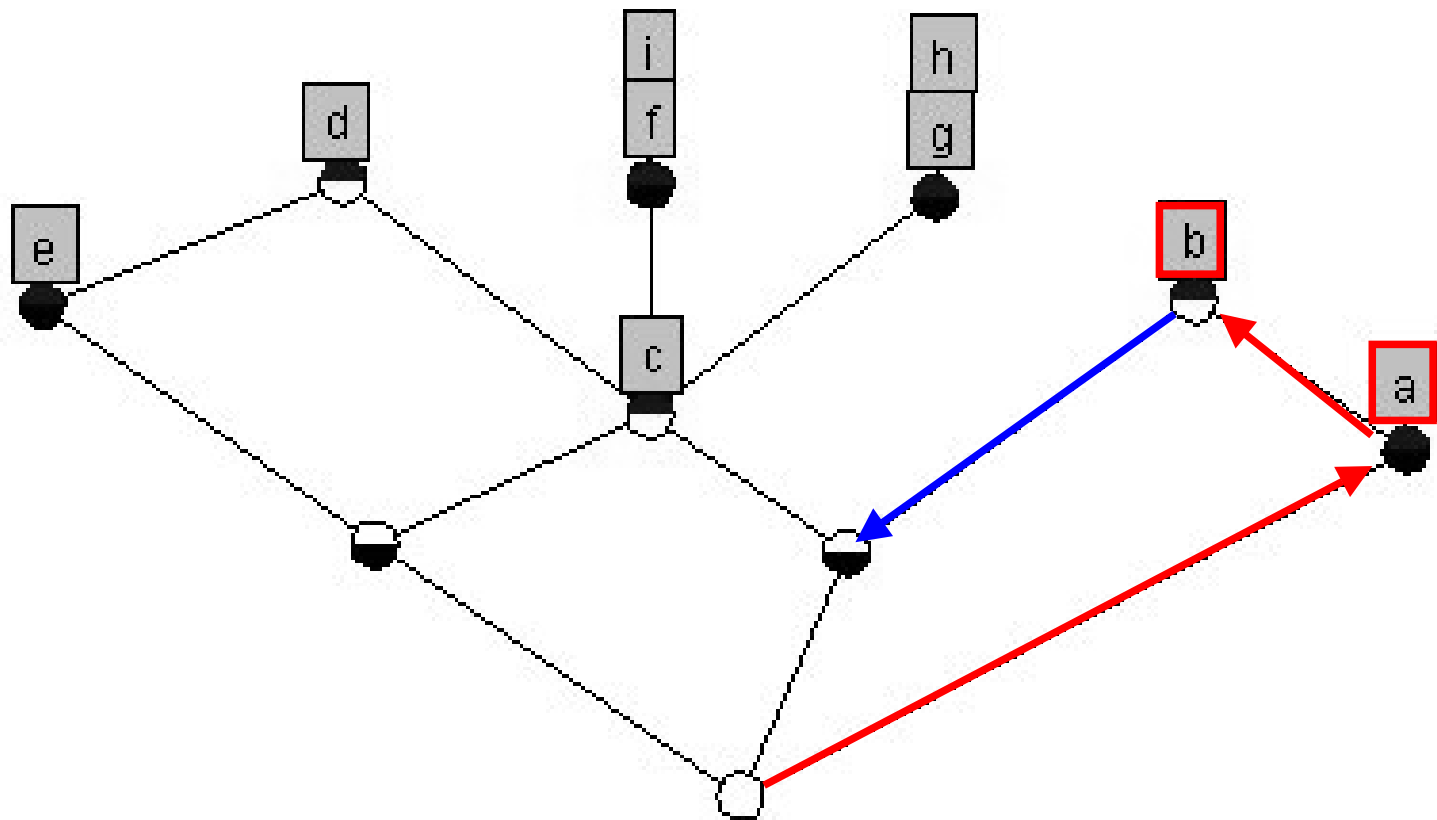


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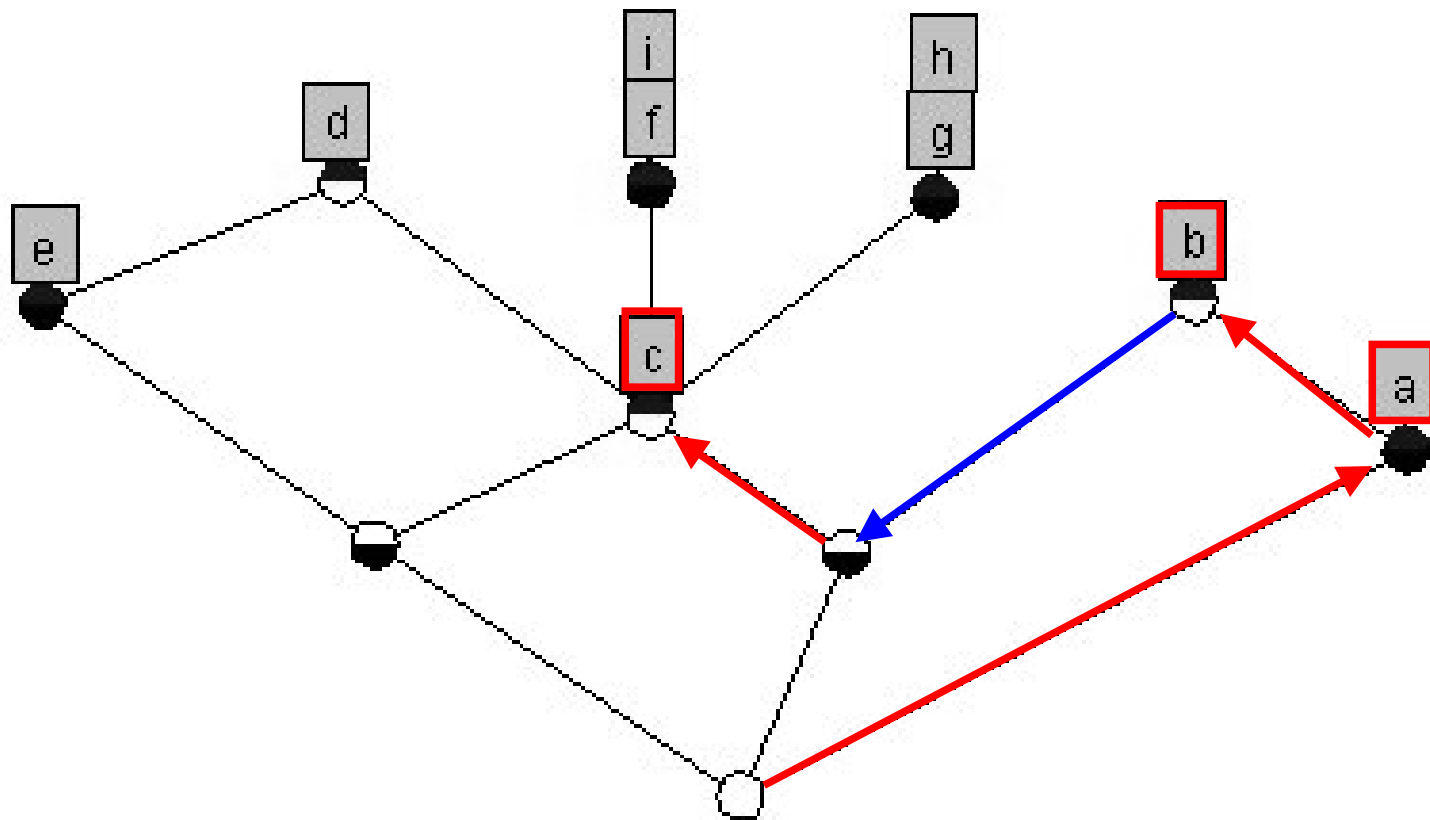
a b

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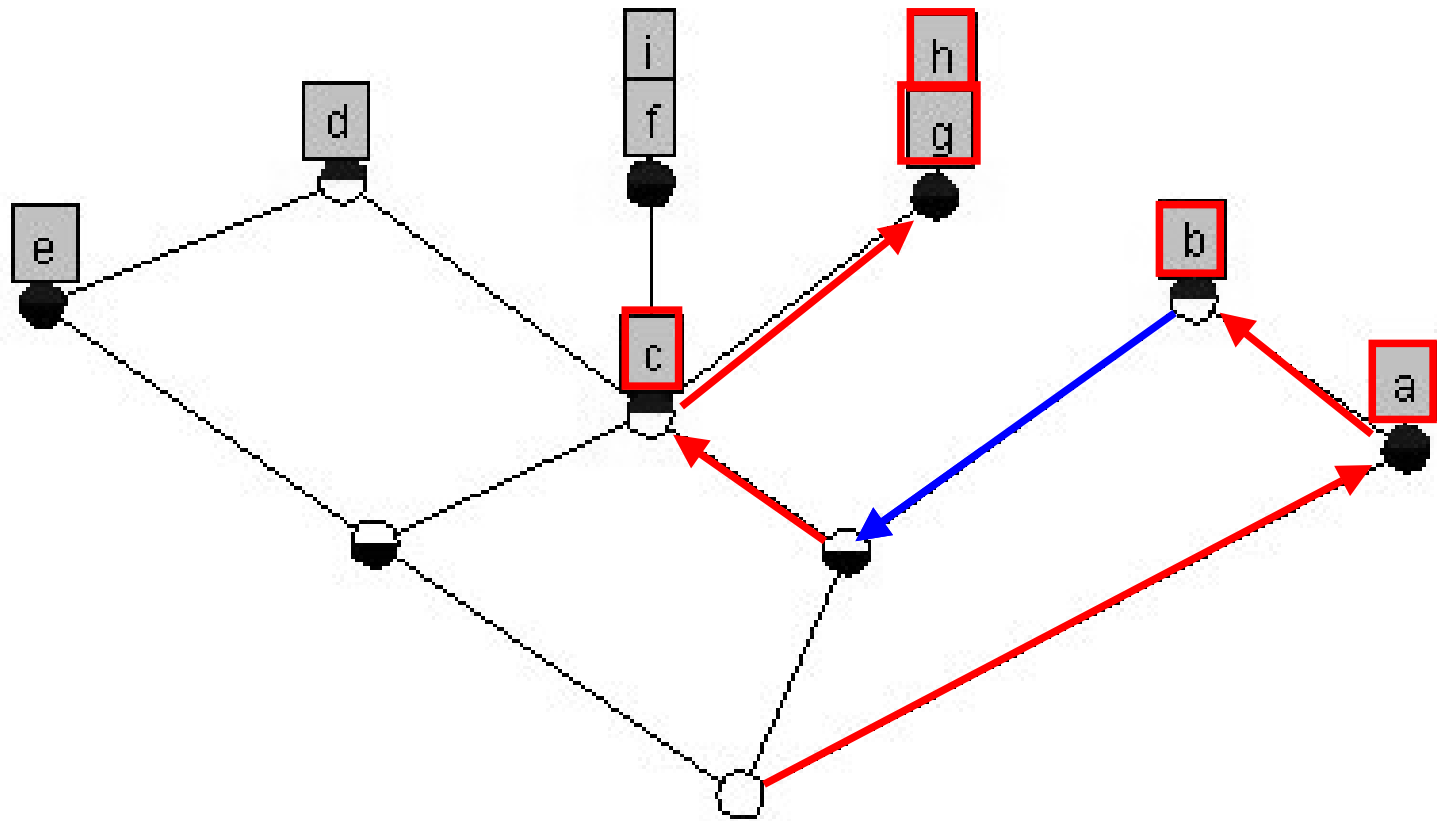
a b  $M_1$

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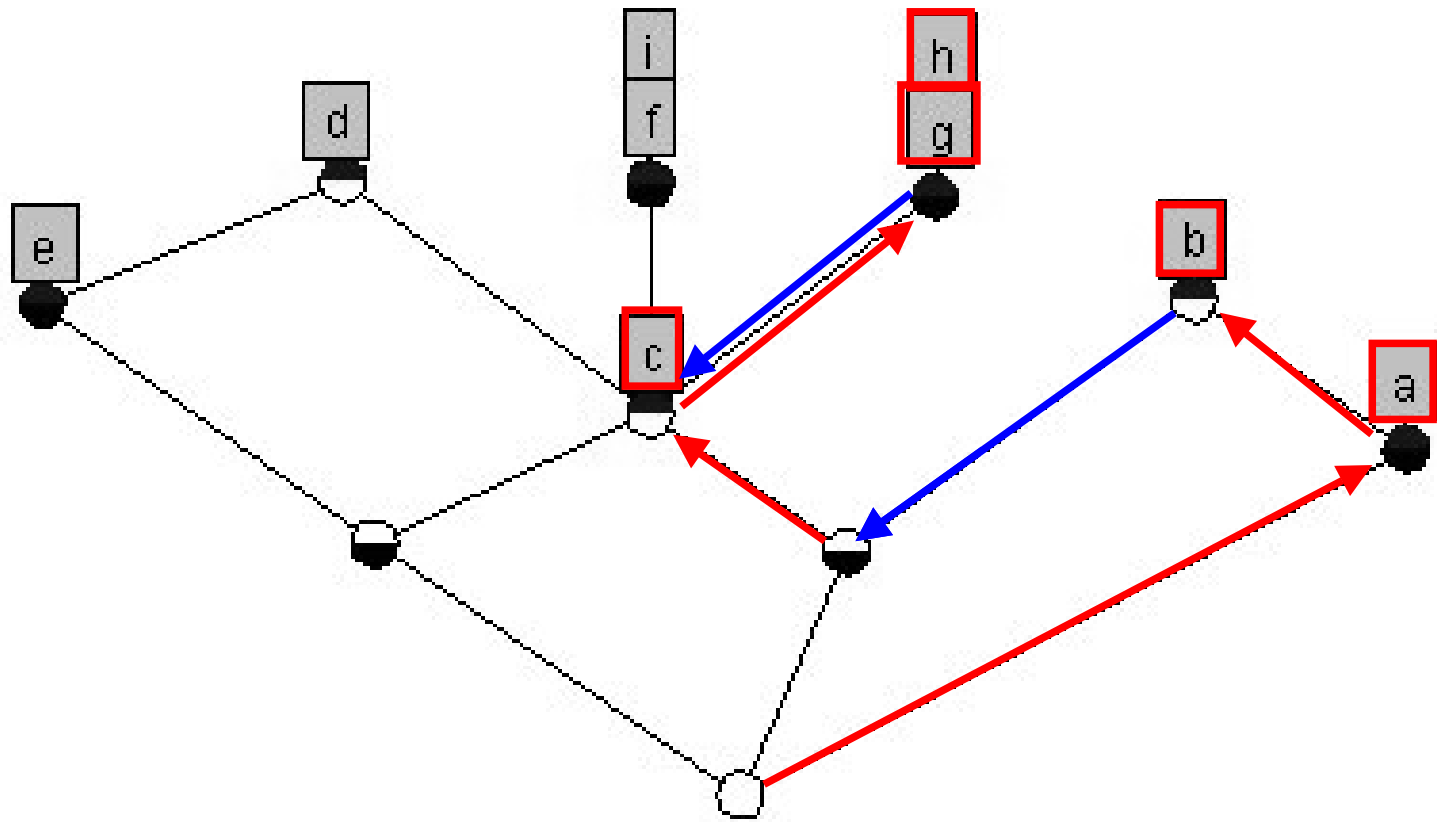
a b  $M_1$  c

# Construction of S-alphabets



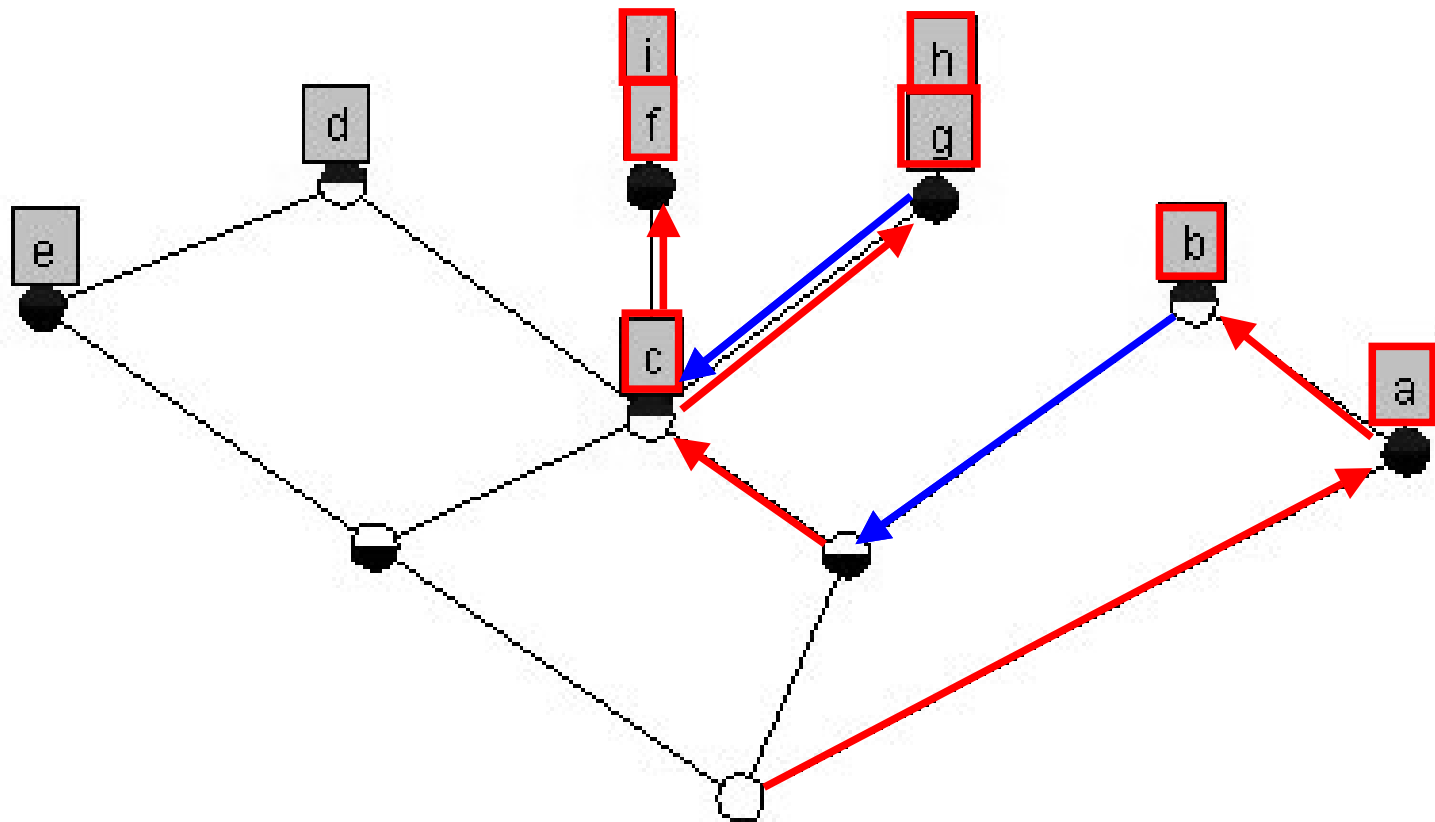
a b  $M_1$  c g h

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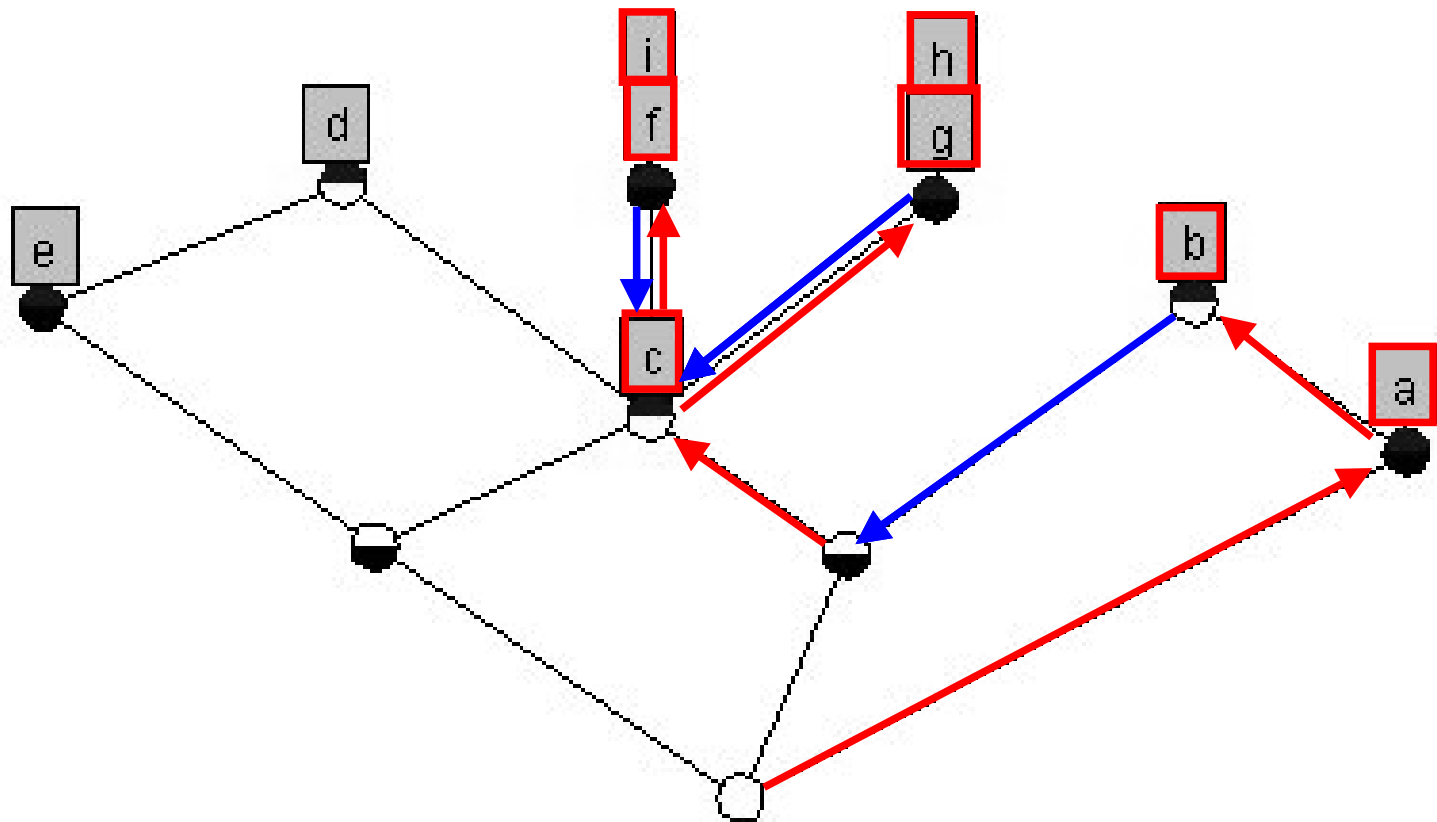
a b  $M_1$  c g h  $M_2$  c

# Construction of S-alphabets



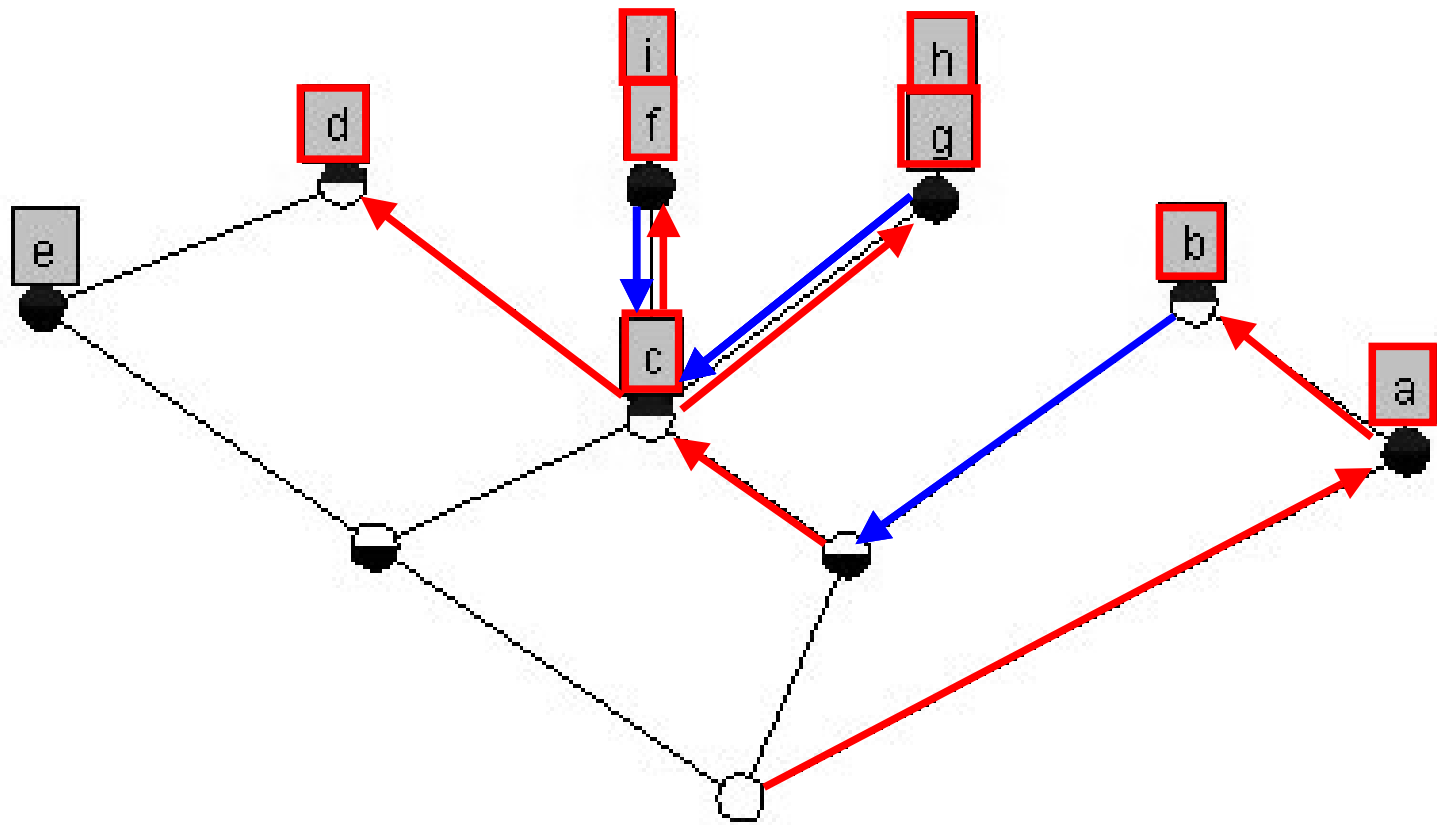
a b  $M_1$  c g h  $M_2$  c i f

# Construction of S-alphabets



a b  $M_1$  c g h  $M_2$  c i f  $M_3$  c

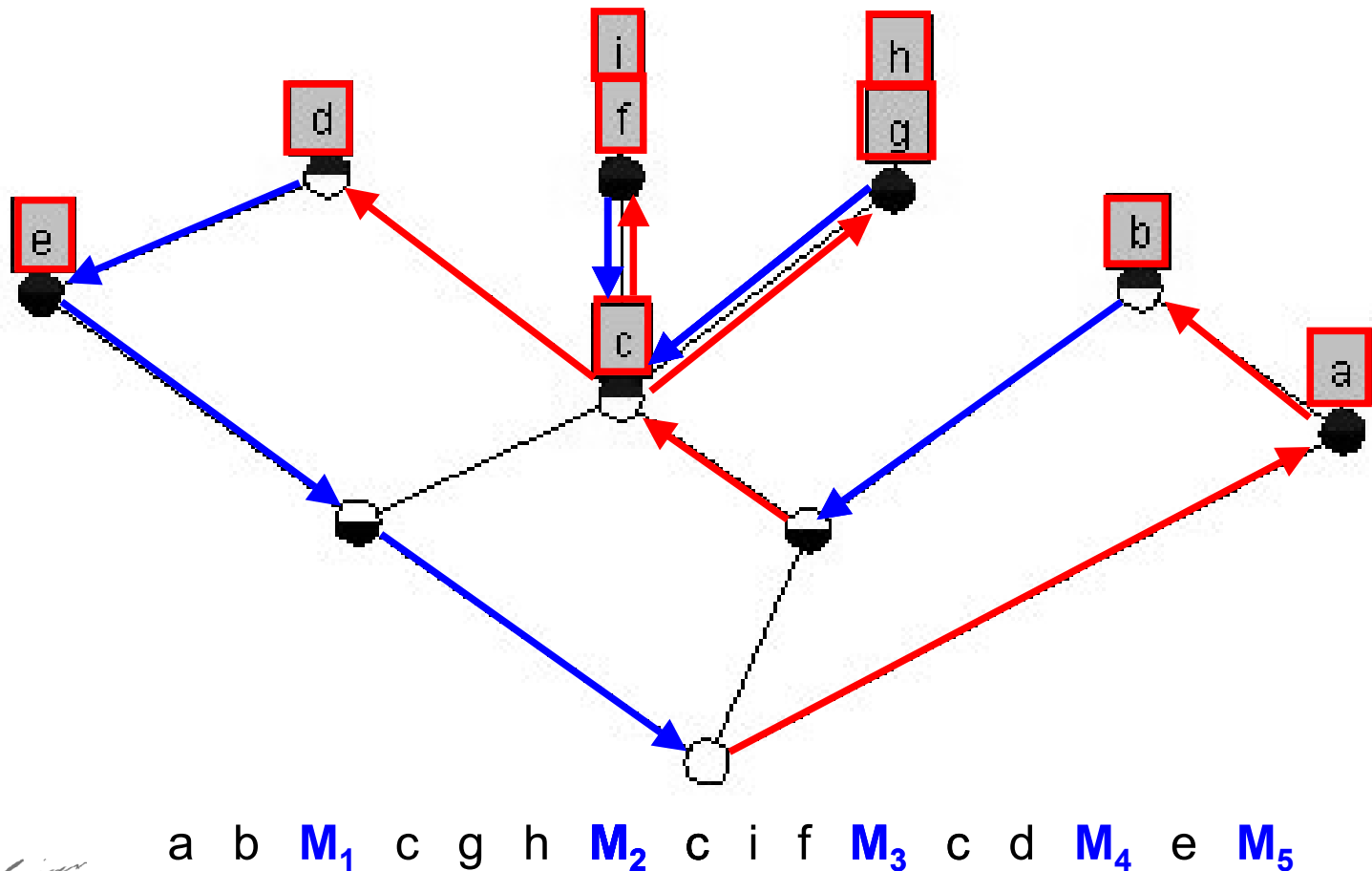
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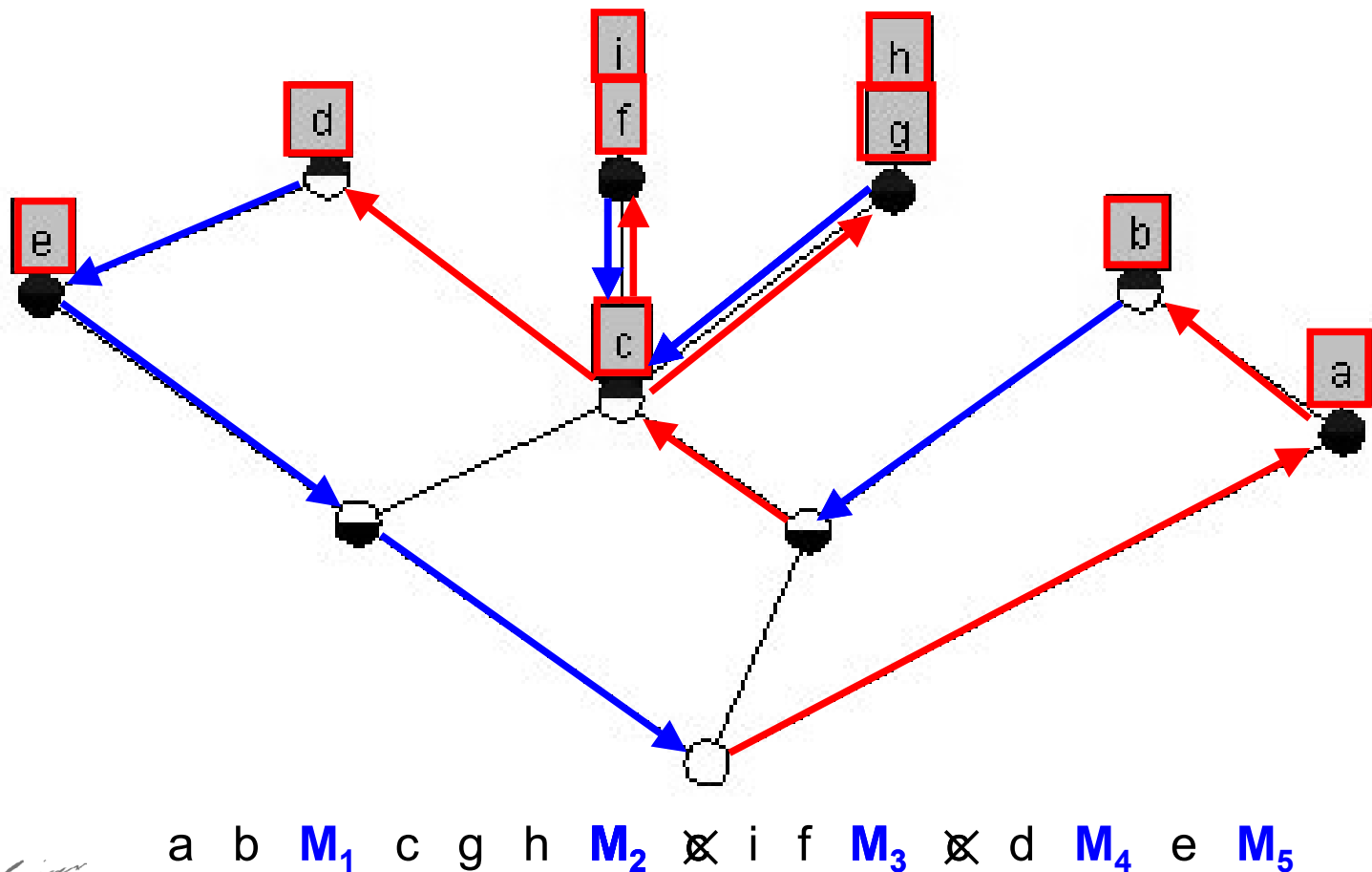
a b  $M_1$  c g h  $M_2$  c i f  $M_3$  c d



# Construction of S-alphabets

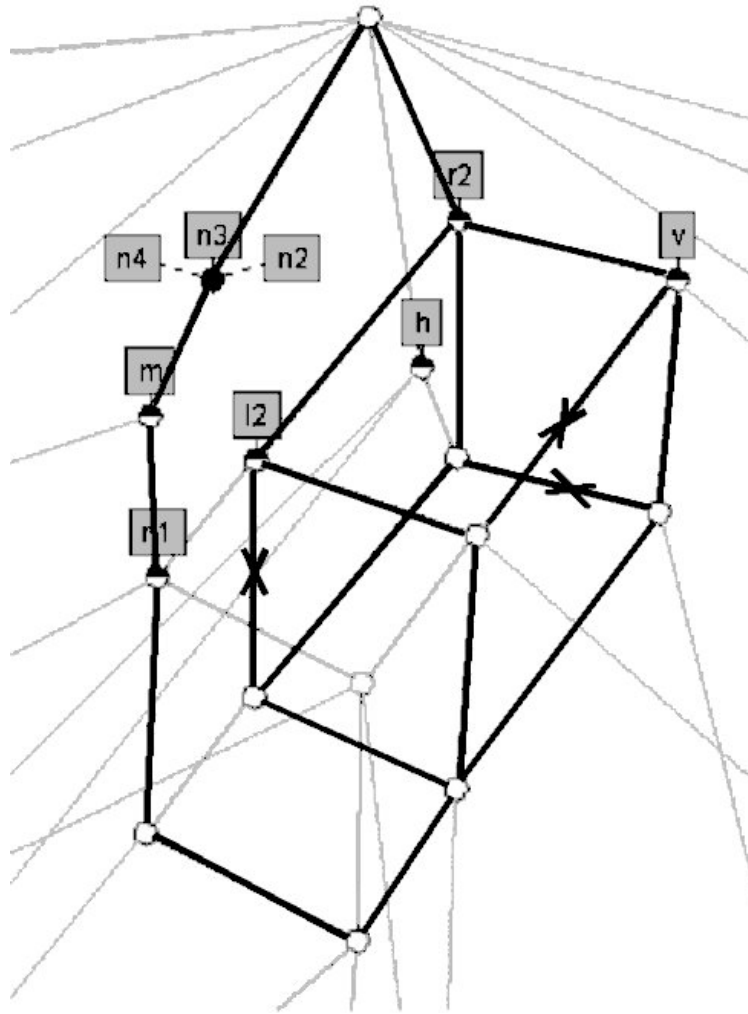


# Construction of S-alphabets









# Pāṇini's Śivasūtras are perfect

