

On the minimality of Pāṇini's *Śivasūtras*

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अइउण्। ऋलृक्। एओङ्। ऐऔच्। हयवरट्।
लण्। अमङणनम्। झभञ्। घढधष्। जबगडदश्।
खफछठथचटतव्। कपय्। शषसर्। हल्।

Phonological Rules

modern notation

A is replaced by B if preceded by C and succeeded by D .

$$A \rightarrow B / C_D$$

example: final devoicing

$$\left[\begin{array}{l} + \text{ consonantal} \\ - \text{ nasal} \\ + \text{ voiced} \end{array} \right] \rightarrow \left[\begin{array}{l} + \text{ consonantal} \\ - \text{ nasal} \\ - \text{ voiced} \end{array} \right] / _ \#$$

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Pāṇini's linear Coding

A + genitive, B + nominative, C + ablative, D + locative.

example

- *sūtra* 6.1.77: *iko yaṇaci* (इको यणचि)
- analysis: $[ik]_{\text{gen}}[yaṇ]_{\text{nom}}[ac]_{\text{loc}}$
- modern notation: $[iK] \rightarrow [yN] / _ [aC]$

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Pāṇini faced the problem of giving a **linear** representation of the **nonlinear** system of sound classes.

A similar problem occurs in . . .

Libraries



Warehouses and stores



Pāṇini's solution: Śivasūtras

1.	a	i	u			N
2.				r	!	K
3.		e	o			Ñ
4.		ai	au			C
5.	h	y	v	r		T
6.					l	N
7.	ñ	m	ṇ	ṇ	n	M
8.	jh	bh				Ñ
9.			gh	ḍh	dh	Ṣ
10.	j	b	g	ḍ	d	Ś
11.	kh	ph	ch	ṭh	th	
			c	ṭ	t	V
12.	k	p				Y
13.		ś	ṣ	s		R
14.	h					L

अइउण्। ऋलृक्।

a·i·uṇ | ṛ·lṛ

एओङ्। ऐऔच्।

e·oṇ | ai·auc

हयवरट्। लण्।

hayavarat | laṇ

ऋमङणनम्। झभञ्।

ṛmaṇaṇanam | jhabhañ

घढधष्। जबगडदश्।

ghaḍhadhaṣ | jabagaḍadaś

खफछठथचटतव्।

khaphachathathacaṭataṇ

कपय्। शषसर्। हल्।

kapay | śaṣasar | hal

Pāṇini's solution: Śivasūtras

1.	a	i	u			Ṇ
2.				r	!	K
3.		e	o			Ṇ
4.		ai	au			C
5.	h	y	v	r		Ṭ
6.					l	Ṇ
7.	ñ	m	ṇ	ṇ	n	M
8.	jh	bh				Ñ
9.			gh	ḍh	dh	Ṣ
10.	j	b	g	ḍ	d	Ś
11.	kh	ph	ch	ṭh	th	
			c	ṭ	t	V
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14.	h					L

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Pāṇini's solution: Śivasūtras

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2.				r	!	K
3.		e	o			Ñ
4.		ai	au			C
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6.					l	N
7.	ñ	m	ṇ	ṇ	n	M
8.	jh	bh				Ñ
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11.	kh	ph	ch	ṭh	th	
			c	ṭ	t	V
12.	k	p				Y
13.		ś	ṣ	s		R
14.	h					L

anubandha

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Pratyāhāras

1.	a	i	u		Ṇ
2.				r	Ḳ
3.		e	o		Ṇ̣
4.		ai	au		C
5.	h	y	v	r	Ṭ

Pratyāhāras

1.	a	i	u		Ṇ
2.				r	Ḳ
3.		e	o		Ṇ
4.		ai	au		C
5.	h	y	v	r	Ṭ

iK

Pratyāhāras

1.	a	i	u		Ṇ
2.				r	Ṛ
3.		e	o		ṅ
4.		ai	au		Ḍ
5.	h	y	v	r	Ṣ

iK = ⟨i, u, r, !⟩

Analysis of iko yaṇaci: [iK] → [yṆ]/_ [aC]

1.	a	i	u			Ṇ
2.				ṛ	ḷ	Ḳ
3.		e	o			Ṇ
4.		ai	au			C
5.	h	y	v	r		Ṭ
6.					l	Ṇ

- [iK] → [yṆ]/_ [aC]
- ⟨i, u, ṛ, ḷ⟩ → ⟨y, v, r, l⟩/_ ⟨a, i, u, ṛ, ḷ, e, o, ai, au⟩

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- [iK] → [yṆ]/_ [aC]
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General problem of S-sortability

Given a set of classes, order the elements of the classes (without duplications) in a linear order (in a list) such that each single class forms a continuous interval with respect to that order.

- The target orders are called **S-orders**
- A set of classes is **S-sortable** if it has an S-order

General problem of Śivasūtra-alphabets (S-alphabets)

Given a set of classes, find an S-order of the elements of the classes. Interrupt this list by markers (*anubandhas*) such that each single class can be denoted by a sound-marker-pair (*pratyāhāra*).

Note that every S-order becomes a Śivasūtra-alphabet (S-alphabet) by adding a marker (*anubandha*) behind each element.

Given the set of classes $\{\{a, b\}, \{a, b, c\}, \{a, b, c, d\}\}$, the order $a b c d$ is one of its S-orders and $a M_1 b M_2 c M_3 d M_4$ is one of its S-alphabets.

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Some more Examples

S-sortable example

The set of classes:

$\{\{d, e\}, \{a, b\}, \{b, c, d, f, g, h, i\}, \{f, i\}, \{c, d, e, f, g, h, i\}, \{g, h\}\}$ is

S-sortable;

one of its S-orders is

a b c g h f i d e

non-S-sortable example

The set of classes:

$\{\{a, b\}, \{b, c\}, \{a, c\}\}$ is not S-sortable.

non-S-sortable example

The set of classes:

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a b c d e or *e d c b a*

Some more Examples

S-sortable example

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a b c g h f i d e

non-S-sortable example

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non-S-sortable example

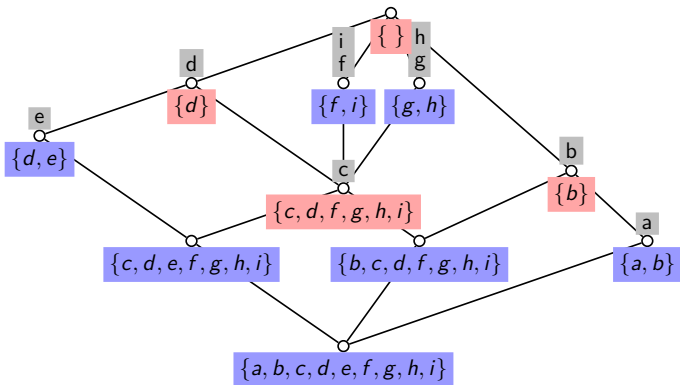
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a b c d e or *e d c b a*

Visualize relations

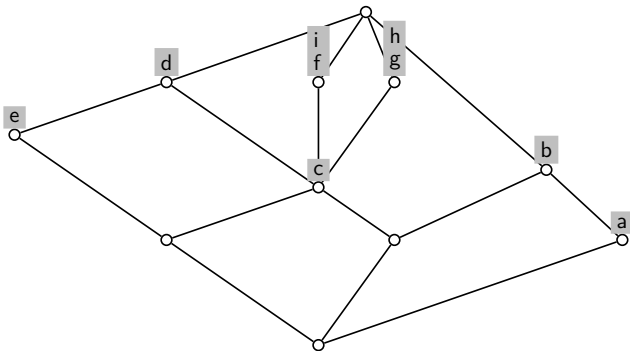
$\{\{d, e\}, \{a, b\}, \{b, c, d, f, g, h, i\}, \{f, i\}, \{c, d, e, f, g, h, i\}, \{g, h\}\}$



concept lattice

Visualize relations

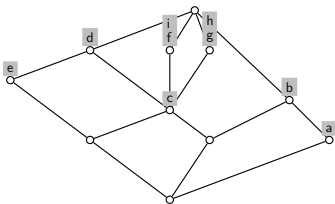
$\{\{d, e\}, \{a, b\}, \{b, c, d, f, g, h, i\}, \{f, i\}, \{c, d, e, f, g, h, i\}, \{g, h\}\}$



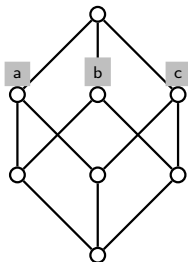
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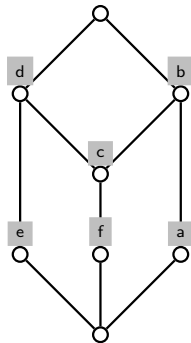
$\{\{d, e\}, \{a, b\}, \{b, c, d, f, g, h, i\}, \{f, i\},$
 $\{c, d, e, f, g, h, i\}, \{g, h\}\}$



$\{\{a, b\}, \{b, c\}, \{a, c\}\}$



$\{\{d, e\}, \{a, b\}, \{b, c, d\},$
 $\{b, c, d, f\}\}$



Getting back to Pāṇini's problem



*a·i·uṇ | ṛ·ḷk | e·oṇ | ai·auc | hayavarat |
laṇ | ṇamaṇaṇanam | jhabhañ | ghaḍhadhaṣ | jabagaḍadaś |
khaphachathathacaṭataṭav | kapay | śaṣasar | hal |*

Q: Are the Śivasūtras minimal (with respect to length)?

What does minimal mean?

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The Śivasūtras are **minimal** if it is **impossible** rearrange the Sanskrit sounds in a new list with *anubandhas* such that

- ① each *pratyāhāra* forms an interval ending before an *anubandha*,
- ② no sound occurs twice
- or** one sound occurs twice but less *anubandhas* are needed.
- ⇒ duplicating a sound is worse than adding *anubandhas*

Common semi-formal argument

Śivasūtras:

aiu **N** ṛ! **K** eo **Ñ** aiau **C** hyvr **T** l **N** ñ m ñ ṇ n **M** jh bh **Ñ** gh ḍh dh **Ṣ**
j b g ḍ d **Ṣ** kh ph ch ṭh th c ṭ t **V** kp **Y** ś ṣ **R** h **L**

$aK = \{a, i, u, r, !\}$, $iK = \{i, u, r, !\}$ and $uK = \{u, r, !\} \Rightarrow a < i < u < r, !$
(taken from Kiparsky 1991) **but:**

$jhL =$

$\{h, s, ṣ, ś, p, k, t, ṭ, c, th, ṭh, ch, ph, kh, d, ḍ, g, b, j, dh, ḍh, gh, bh, jh\}$

$jhR =$

$\{s, ṣ, ś, p, k, t, ṭ, c, th, ṭh, ch, ph, kh, d, ḍ, g, b, j, dh, ḍh, gh, bh, jh\}$

$jhY = \{p, k, t, ṭ, c, th, ṭh, ch, ph, kh, d, ḍ, g, b, j, dh, ḍh, gh, bh, jh\}$

$jhṢ = \{d, ḍ, g, b, j, dh, ḍh, gh, bh, jh\}$ and

$jhṢ = \{dh, ḍh, gh, bh, jh\}$

$\Rightarrow h < s, ṣ, ś < p, k, t, ṭ, c, th, ṭh, ch, ph, kh, d < ḍ, g, b, j <$

$dh, ḍh, gh, bh, jh$

Common semi-formal argument

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aiu **N** ṛ! **K** eo **Ñ** aiau **C** hyvr **T** l **N** ñ m ñ ṇ n **M** jh bh **Ñ** gh ḍh dh **Ṣ**
j b g ḍ d **Ṣ** kh ph ch ṭh th c ṭ t **V** kp **Y** ś ṣ s **R** h **L**

$aK = \{a, i, u, \dot{r}, !\}$, $iK = \{i, u, \dot{r}, !\}$ and $uK = \{u, \dot{r}, !\} \Rightarrow a < i < u < \dot{r}, !$
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$jhL =$

$\{h, s, \dot{s}, \acute{s}, p, k, t, \dot{t}, c, th, \dot{t}h, ch, ph, kh, d, \dot{d}, g, b, j, dh, \dot{d}h, gh, bh, jh\}$

$jhR =$

$\{s, \dot{s}, \acute{s}, p, k, t, \dot{t}, c, th, \dot{t}h, ch, ph, kh, d, \dot{d}, g, b, j, dh, \dot{d}h, gh, bh, jh\}$

$jhY = \{p, k, t, \dot{t}, c, th, \dot{t}h, ch, ph, kh, d, \dot{d}, g, b, j, dh, \dot{d}h, gh, bh, jh\}$

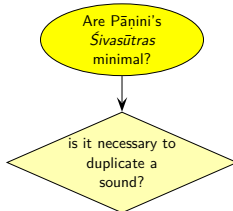
$jh\acute{S} = \{d, \dot{d}, g, b, j, dh, \dot{d}h, gh, bh, jh\}$ and

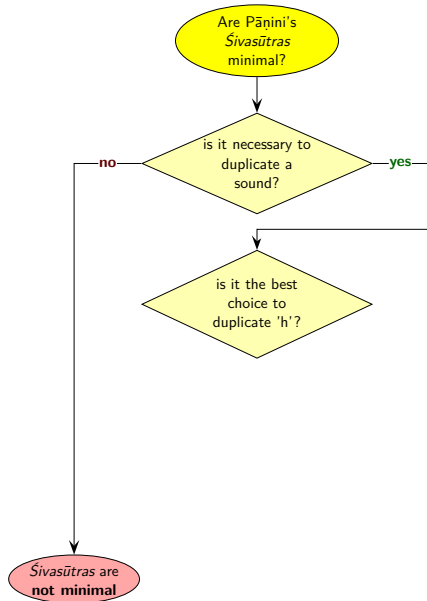
$jh\dot{S} = \{dh, \dot{d}h, gh, bh, jh\}$

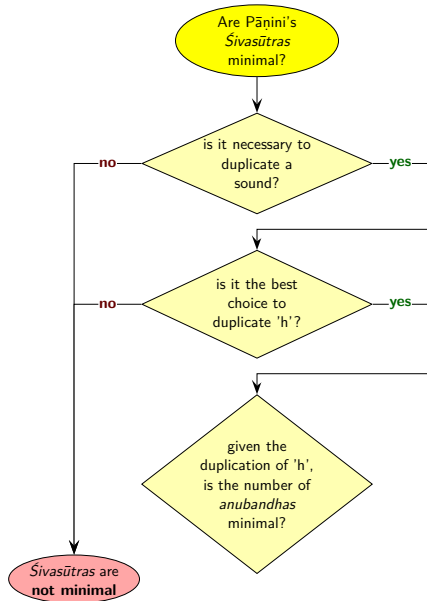
$\Rightarrow h < s, \dot{s}, \acute{s} < p, k, t, \dot{t}, c, th, \dot{t}h, ch, ph, kh, d < \dot{d}, g, b, j <$

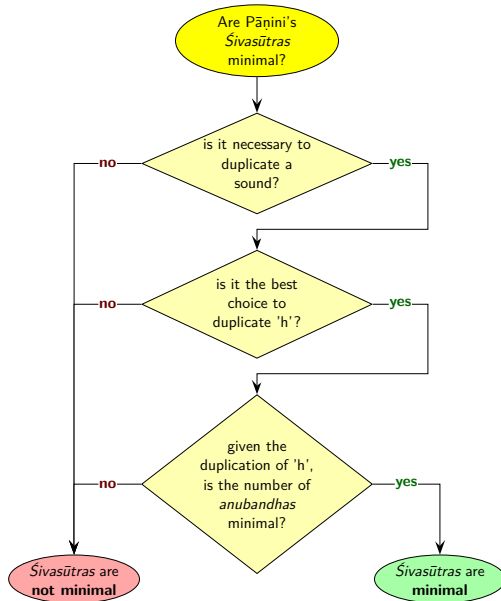
$dh, \dot{d}h, gh, bh, jh$

Are Pāṇini's
Śivasūtras
minimal?





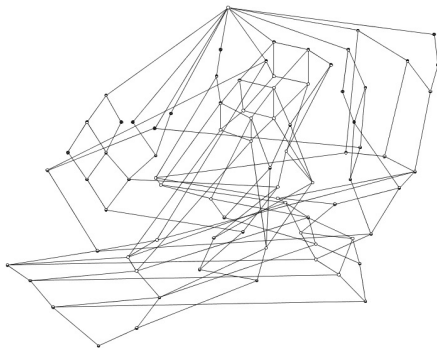
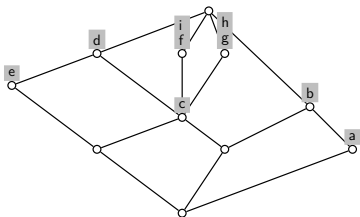




Is it necessary to duplicate a sound?

Main theorem on S-sortability (part 1a)

If a set of classes is S-sortable, then its concept lattice is Hasse-planar.

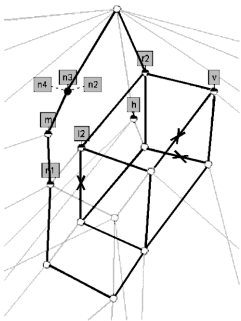


graph of the concept lattice of Pāṇini's
pratyāhāras

No S-alphabet without duplications for Pāṇini's *pratyāhāras*

Criterion of Kuratowski

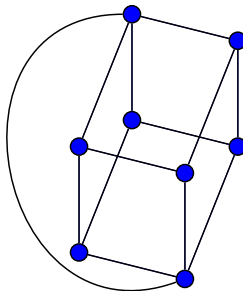
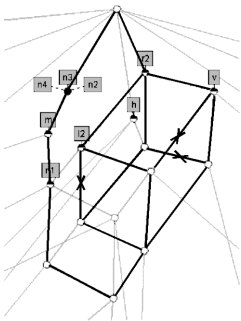
A graph which has the graph  as a minor is not planar.



No S-alphabet without duplications for Pāṇini's *pratyāhāras*

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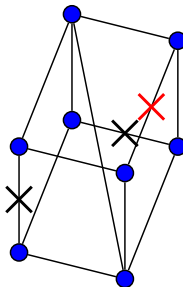
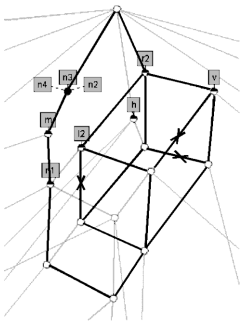
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No S-alphabet without duplications for Pāṇini's *pratyāhāras*

Criterion of Kuratowski

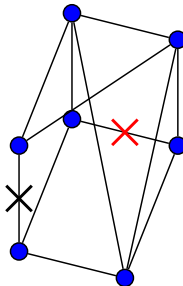
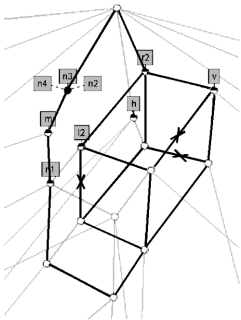
A graph which has the graph  as a minor is not planar.



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Criterion of Kuratowski

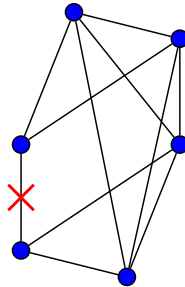
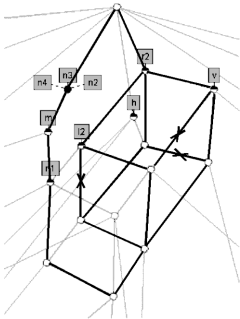
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Criterion of Kuratowski

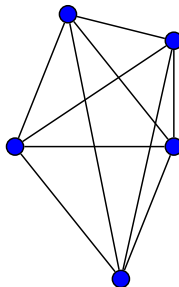
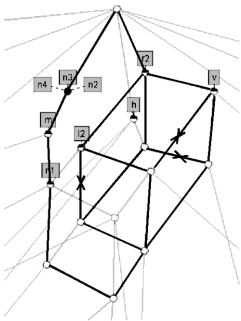
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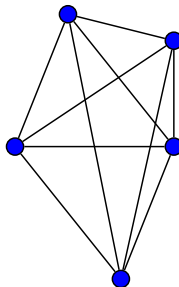
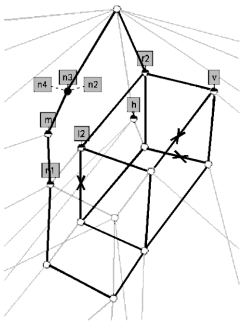
A graph which has the graph  as a minor is not planar.



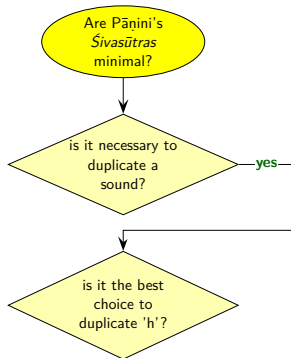
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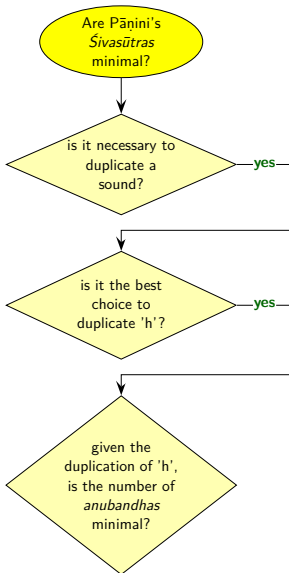
Criterion of Kuratowski

A graph which has the graph  as a minor is not planar.

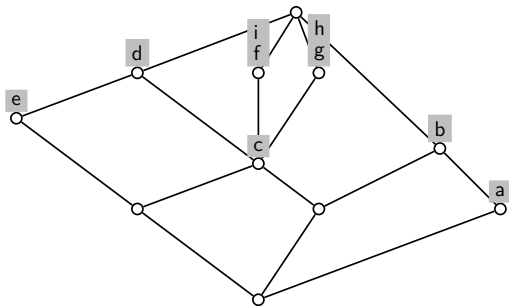


There is no S-alphabet for the set of classes given by Pāṇini's *pratyāhāras* **without** duplicated elements!





S-alphabets with a minimal number of markers

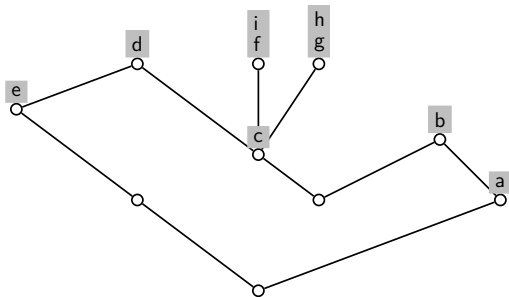


procedure

Start with the empty sequence and choose a walk through the S-graph:

- While moving upwards do nothing.
- While moving downwards along an edge add a new marker to the sequence unless its last element is already a marker.
- If a labeled node is reached, add the labels in arbitrary order to the sequence, unless it has been added before.

S-alphabets with a minimal number of markers

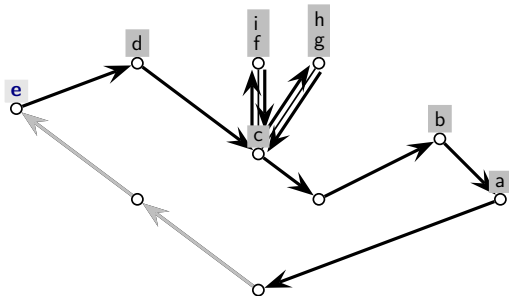


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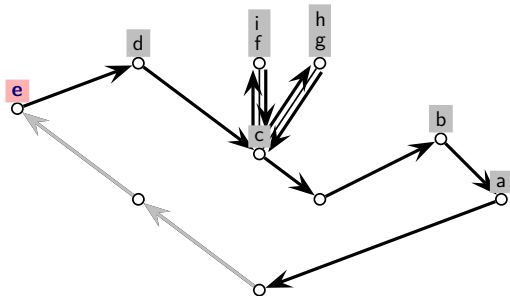


procedure

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- If a sound is reached, add the sound to the sequence, unless it has been added before.

S-alphabets with a minimal number of markers



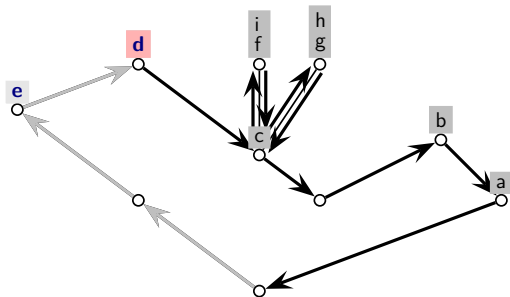
e

procedure

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S-alphabets with a minimal number of markers



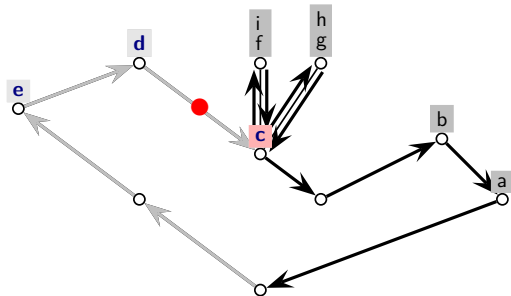
ed

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S-alphabets with a minimal number of markers



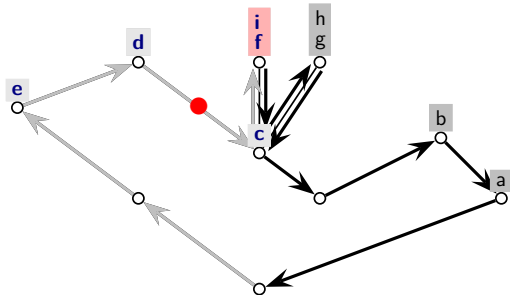
edM_1c

procedure

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S-alphabets with a minimal number of markers



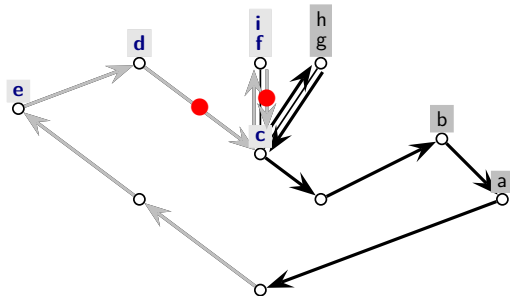
edM₁cfi

procedure

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S-alphabets with a minimal number of markers



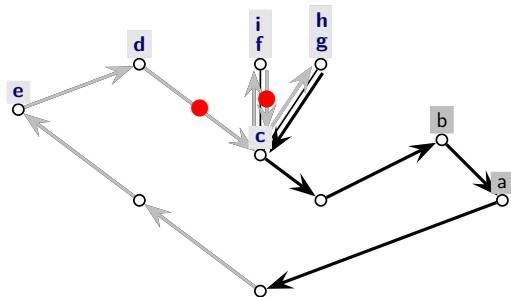
edM_1cfiM_2

procedure

Start with the empty sequence and choose a walk through the S-graph:

- While moving upwards do nothing.
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S-alphabets with a minimal number of markers



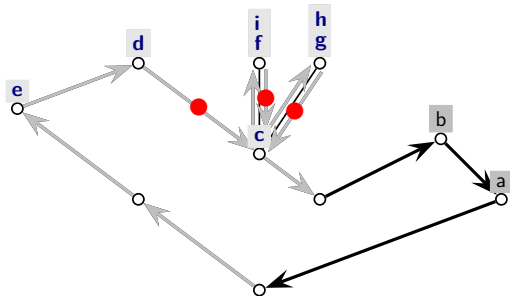
edM₁cfiM₂gh

procedure

Start with the empty sequence and choose a walk through the S-graph:

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S-alphabets with a minimal number of markers



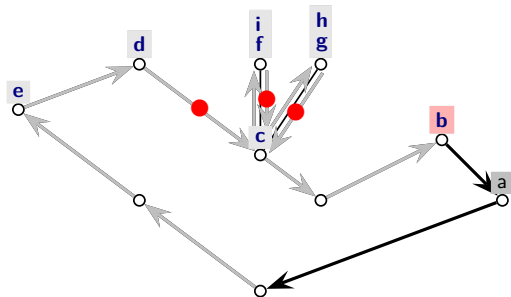
$edM_1cfiM_2ghM_3$

procedure

Start with the empty sequence and choose a walk through the S-graph:

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S-alphabets with a minimal number of markers



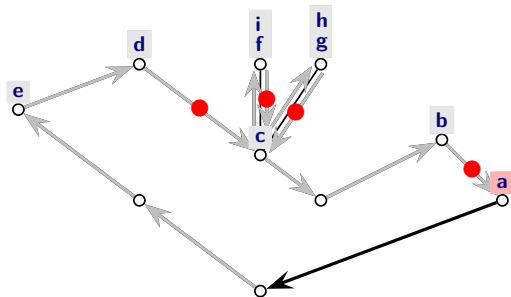
$edM_1cfiM_2ghM_3b$

procedure

Start with the empty sequence and choose a walk through the S-graph:

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S-alphabets with a minimal number of markers



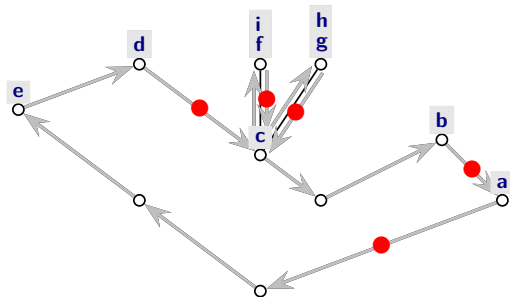
$edM_1cfiM_2ghM_3bM_4a$

procedure

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S-alphabets with a minimal number of markers



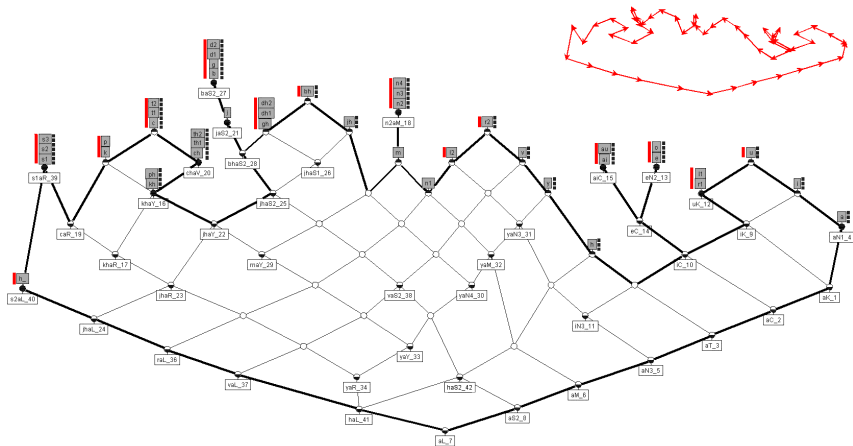
$edM_1cfiM_2ghM_3bM_4aM_5$

procedure

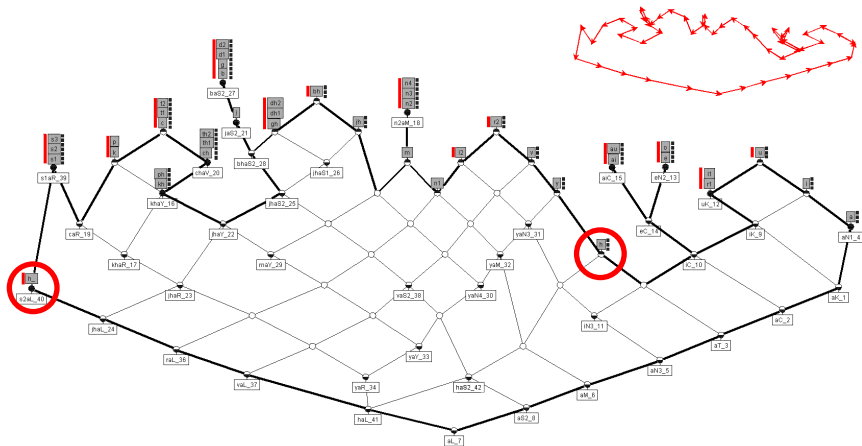
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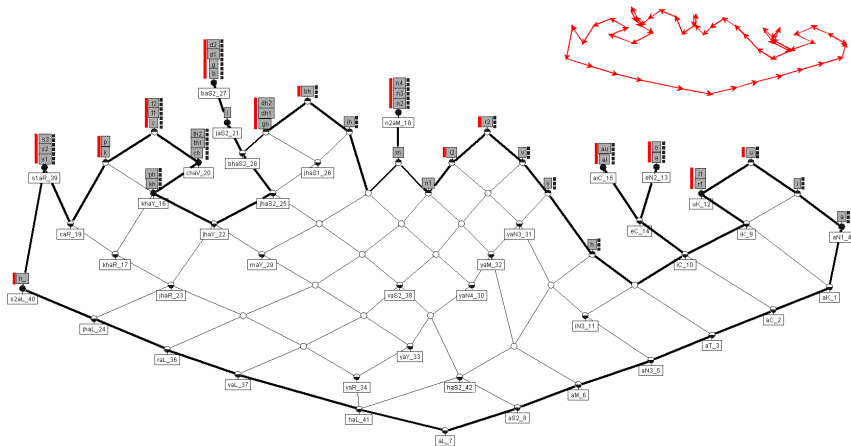
Concept lattice of Pāṇini's *pratyāhāras* with duplicated *h*



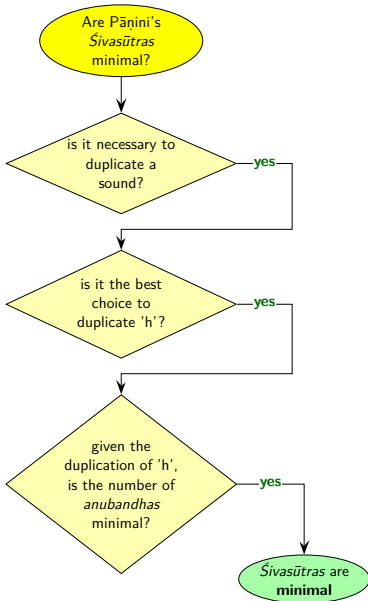
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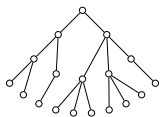


With the Śivasūtras Pāṇini has chosen one out of nearly 12 million minimal S-alphabets!

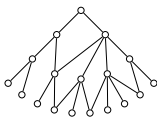


Transfer

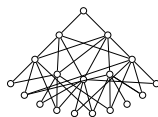
- For physical objects ,duplicating' means ,adding copies'
- Adding copies is annoying but often not impossible
- Ordering objects in an S-order may
 - improve user-friendliness
 - save time
 - save space
 - simplify visual representations of classifications



tree



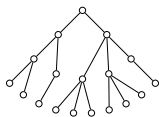
S-sortable



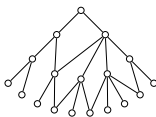
general hierarchy

Transfer

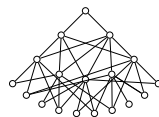
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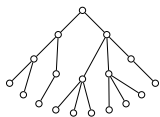
S-sortable



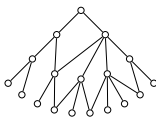
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Transfer

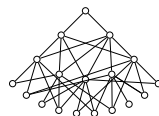
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tree



S-sortable



general hierarchy

Transfer

Objects in libraries, ware-houses, and stores are only *nearly* linearly arranged:

⇒ Second (and third) dimension can be used in order to avoid duplications



Open problems

What explains the actual structure of the *Śivasūtras*?

- **principle of homorganic continuity** (Staal, 1962)
- **principle of historic continuity** (Cardona, 1969)
- **principle of economy and logic of the special case and the general case** (Kiparsky 1991) or **Pāṇini's razor** (Kiparsky 2007)

The presented approach cannot give an answer to this question

The story is much more intricate

- We have **neither** shown that Pāṇini's technique for the representation of sound classes is optimal
- **nor** that he has used his technique in an optimal way.

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Literature

-  Cardona, G. (1969), Studies in Indian grammarians I: The method of description reflected in the Śiva-Sūtras. In: Transactions of the American Philosophical Society, 59(1):3–48.
-  Kiparsky, P. (1991), Economy and the construction of the Śivasūtras. In: M. M. Deshpande & S. Bhate (eds.), *Pāṇinian Studies*, Michigan: Ann Arbor.
-  Petersen, W. (2008), Zur Minimalität von Pāṇinis *Śivasūtras* – Eine Untersuchung mit Mitteln der Formalen Begriffsanalyse. PhD thesis, university of Düsseldorf.
-  Petersen, W. (2009), On the Construction of Sivasutra-Alphabets. In: A. Kulkarni and G. Huet (eds.): *Sanskrit Computational Linguistics*. LNCS 5406, Springer.
-  Staal, F. (1962), A Method of Linguistic Description. *Language* **38**, 1-10.

Origin of Pictures

- libraries (left):
<http://www.meduniwien.ac.at/medizinischepsychologie/bibliothek.htm>
- libraries (middle): <http://www.math-nat.de/aktuelles/allgemein.htm>
- libraries (right):
<http://www.geschichte.mpg.de/deutsch/bibliothek.html>
- warehouses:
http://www.metrogroup.de/servlet/PB/menu/1114920_l1/index.html
- stores: <http://www.einkaufsparadies-schmidt.de/01bilder01/>