# Parsing Beyond Context-Free Grammars: Natural Languages are not Context-Free

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### **Overview**

- 1 CFG and Natural Languages
- 2 Cross-serial Dependencies
- 3 Swiss German is not Context-Free
- 4 Beyond CFG: TAG and LCFRS

### **CFG** and Natural Languages

- For a long time there has been a debate about whether CFGs are sufficiently powerful to describe natural languages. Several approaches have used CFGs, oftentimes enriched with some additional mechanism of transformation [Cho56] or with features [GKPS85] for natural languages.
- In the 80's Stuart Shieber was able to prove in [Shi85] that there
  are natural languages that cannot be generated by a CFG. Before
  that, [BKPZ82] made already a similar argument but their proof
  is based on the tree structures obtained with CFGs while Shieber
  argues on the basis of weak generative capacity, i.e., of the string
  languages.
- The phenomena considered in both papers are *cross-serial* dependencies.

# **Cross-serial Dependencies (1)**

Cross-serial dependencies in Dutch [BKPZ82]:

```
    ... dat Jan de kinderen zag zwemmen
    ... that Jan the children saw swim
    ... that Jan saw the children swim'
```

The colours mark the dependencies between the two verbs and the two NPs: *the children* is an argument of *swim* while *Jan* is an argument of *saw*. The dependency links are in a crossing configuration.

# **Cross-serial Dependencies (2)**

#### This phenomenon can be iterated:

- (2) ... dat Jan Piet de kinderen zag helpen zwemmen... that Jan Piet the children saw help swim' ... that Jan saw Piet help the children swim'
- (3) ... dat Jan Piet Marie de kinderen zag helpen leren zwemmen
  ... that Jan Piet Marie the children saw help teach swim
  '... that Jan saw Piet help Marie teach the children to swim'

# Cross-serial Dependencies (3)

- In principle, an unbounded number of crossed dependencies is possible.
- However, except for the first and last verb any permutation of the NPs and the verbs is grammatical as well (even though with a completely different dependency structure since the dependencies are always cross-serial).
- Therefore, the dependencies are not visible on the strings and the string language of Dutch cross-serial dependencies amounts roughly to  $\{n^k v^k \mid k > 0\}$  which is a context-free language.

This is different for Swiss German because Swiss German has case marking.

# **Cross-serial Dependencies (4)**

### Cross-serial dependencies in Swiss German [Shi85]:

- (4) ... das mer em Hans es huus hälfed aastriiche ... that we  $\operatorname{Hans}_{Dat}$  house<sub>Acc</sub> helped paint ... that we helped Hans paint the house'
- (5)... das mer d'chind em Hans es huus lönd hälfe ... that we the children<sub>Acc</sub>  $Hans_{Dat}$  house<sub>Acc</sub> let aastriiche paint ... that we let the children help Hans paint the house'

#### Swiss German

- uses case marking
- and displays cross-serial dependencies.

### Swiss German is not Context-Free (1)

#### **Proposition 1**

The language L of Swiss German is not context-free [Shi85].

The argumentation of the proof goes as follows:

- We assume that I is context-free.
- Then the intersection of a regular language with the image of L under a homomorphism must be context-free as well.
- We find a particular homomorphism and a regular language such that the result obtained in this way is a non context-free language.
- This is a contradiction to our assumption and, consequently, the assumption does not hold.

## Swiss German is not Context-Free (2)

Consider sentences of the following form:

```
(6) ... das mer d'chind em Hans es huus haend
... that we the children<sub>Acc</sub> Hans<sub>Dat</sub> house<sub>Acc</sub> have
wele laa hälfe aastriiche
wanted let help paint
' ... that we let the children help Hans paint the house'
```

The NP verb pairs d'chind laa and em Hans hälfe both can be iterated.

### Swiss German is not Context-Free (3)

With an additional embedding under Jan säit ('Jan says') we obtain constructions of the form

(7) Jan säit das mer (d'chind)<sup> $i_1$ </sup> (em Hans)<sup> $j_1$ </sup> (d'chind)<sup> $i_2$ </sup> (em Hans)<sup> $j_2$ </sup> ... es huus haend wele (laa)<sup> $i_1$ </sup> (hälfe)<sup> $j_1$ </sup> (laa)<sup> $i_2$ </sup> (hälfe)<sup> $j_2$ </sup> ... aastriiche

#### where

- the number of accusative NPs d'chind must equal the number of verbs (here laa) selecting for an accusative,
- the number of dative NPs em Hans must equal the number of verbs (here hälfe) selecting for a dative object, and
- the order of NPs and verbs must be the same in the sense that if all accusative NPs precede all dative NPs, then all verbs selecting an accusative must precede all verbs selecting a dative.

### Swiss German is not Context-Free (4)

The following homomorphism f separates the iterated noun phrases and verbs in these examples from the surrounding material:

```
f(\text{``d'chind''}) = a
f(\text{``em Hans''}) = b
f(\text{``laa''}) = c
f(\text{``halfe''}) = d
f(\text{``Jan sait das mer''}) = w
f(\text{``es huus haend wele''}) = x
f(\text{``aastriiche''}) = y
f(s) = z otherwise
```

### Swiss German is not Context-Free (5)

The images of the constructions we are interested in under f are of the form  $wv_1xv_2y$  where  $v_1$  contains as and bs and  $v_2$  contains cs and ds and if the kth element in  $v_1$  is an a (a b resp.), then the kth element in  $v_2$  is a c (a d resp.). All other sentences have a c somewhere in their image under f.

To make sure we concentrate only on the constructions of the described form and only on constructions where the accusative NPs precede the dative NPs, we intersect f(L) with the regular language  $wa^*b^*xc^*d^*y$ .

$$L' = f(L) \cap wa^*b^*xc^*d^*y = \{wa^ib^jxc^id^jy \mid i, j \ge 0\}$$

### Swiss German is not Context-Free (6)

If L is context-free, then L' must be context-free as well.

• Then the image of L' under a homomorphism f' with  $f'(w) = f'(x) = f'(y) = \varepsilon$ , f'(a) = a, f'(b) = b, f'(c) = c, f'(d) = d is also context-free. This image is

$$f'(L') = L'' = \{a^i b^j c^i d^j \mid i, j \ge 0\}$$

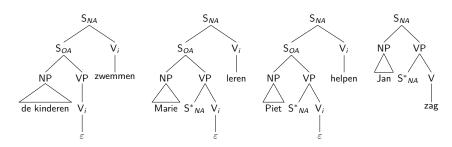
• Consequently, L'' satisfies the pumping lemma for context-free languages. Inspecting the word  $a^k b^k c^k d^k$  where k is the constant from the pumping lemma, this can be shown to lead to a contradiction.

Consequently, L'' is not context-free, and neither are L' and L.

### TAG and Cross-serial Dependencies

TAG for Dutch cross-serial dependencies:

(8) ... dat Jan Piet Marie de kinderen zag helpen leren zwemmen
... that Jan Piet Marie the children saw help teach swim
' ... that Jan saw Piet help Marie teach the children to swim'



the NPs are of course be added by substitution

## **LCFRS** and Cross-serial Dependencies

LCFRS for Dutch cross-serial dependencies:

```
\begin{array}{cccccc} \mathsf{S}(X \ Y \ \mathsf{zag} \ Z) & \to & \mathsf{NP}(X) \ \mathsf{VP}(Y,Z) \\ \mathsf{VP}(X \ Y, \mathsf{leren} \ Z) & \to & \mathsf{NP}(X) \ \mathsf{VP}(Y,Z) \\ \mathsf{VP}(X \ Y, \mathsf{helpen} \ Z) & \to & \mathsf{NP}(X) \ \mathsf{VP}(Y,Z) \\ \mathsf{VP}(X, \mathsf{zwemmen}) & \to & \mathsf{NP}(X) \\ & & \mathsf{NP}(Jan) & \to & \varepsilon \\ & & \mathsf{NP}(Marie) & \to & \varepsilon \\ & & \mathsf{NP}(Piet) & \to & \varepsilon \\ & & \mathsf{NP}(\mathsf{de} \ \mathsf{kinderen}) & \to & \varepsilon \end{array}
```

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