# Parsing Beyond CFG <br> Homework 2: TAG 

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## Question 1

Consider a TAG $G$ with start symbol $S$ and only the following two trees:
$\alpha$ :
${ }_{\varepsilon}^{S}$ $\beta$ :


1. What is the string language generated by this TAG?
2. Now consider a TAG $G^{\prime}$ that contains the same trees but without adjunction constraints. Does it generate the same string language as G? If not, give an example of a string that is only in one of the two languages.
3. What is the string language of $G^{\prime}$ ?

Solution:

1. $L(G)=\left\{a^{n} b^{n} c^{n} d^{n} \mid n \geq 0\right\}$
2. By using $\beta$ twice, adjoining each time to the root (i.e., $\beta$ adjoins to the root of $\beta$, and this derived auxiliary tree adjoins at the root of $\alpha$, we obtain a tree for the string $a b a b c d c d \notin L(G)$.
3. Clearly, if the leftmost terminal is a $a$ (resp. a $b$ ), the rightmost is a $d$ (resp. a $c$ ). The same holds for the next pair of terminals when moving towards the middle of a word, and so on. Furthermore, whenever adding a pair $a$ and $b$ to the first half of a word, the $a$ precedes the $b$. Otherwise, the new pair can be added in any position. This means that, in the first half of the word, $w_{1}$, we have $\left|w_{1}\right|_{a}=\left|w_{1}\right|_{b}$ and for any two $v_{1}, v_{2}$ with $w_{1}=v_{1} v_{2},\left|v_{1}\right|_{a} \geq\left|v_{2}\right|_{a}$ holds.
$L\left(G^{\prime}\right)=\left\{w_{1} w_{2}\left|w_{1} \in\{a, b\}^{*},\left|w_{1}\right|_{a}=\left|w_{1}\right|_{b}\right.\right.$, for any $v_{1}, v_{2}$ with $w_{1}=v_{1} v_{2},\left|v_{1}\right|_{a} \geq\left|v_{2}\right|_{a}$ holds, and, furthermore, $w_{2}$ is the image of $w_{1}^{R}$ under a homomorphism $f$ with $f(a)=$ $d, f(b)=c\}$

## Question 2

Now consider the following elementary trees:
$\alpha$ :
$\beta_{1}$ :

$\beta_{2}$ :


Add adjunction constraints of the form $N A$ (stands for $f_{S A}(v)=\emptyset$ ), $O A(X)$ (stands for $f_{O A}(v)=1$ and $f_{S A}(v)=X$ ) or $S A(X)$ (stands for $f_{O A}(v)=0$ and $f_{S A}(v)=X$ ) to the trees so that the TAG generates $\left\{a^{n} b^{n} c^{n} d^{m} e^{m} f^{m} \mid n, m \geq 1\right\}$.

Solution
$\alpha$ :

$\beta_{1}$ :

$\beta_{2}$ :


## Question 3

Consider the derivation for "Bill John claims Mary likes" on slide 25, ignoring the subscript $i$.

1. Give the resulting derived tree.
2. Give the resulting derivation tree. (As names for the trees, you can use their lexical anchors.)

Solution
1.

2.


